Algorithms for Model Checking (2IW55)

Lecture 6
Parity games

Background material: Chapter 3 of
J.J.A. Keiren, *An experimental study of algorithms and optimisations for parity games, with an application to Boolean Equation Systems*, MSc thesis, 2009

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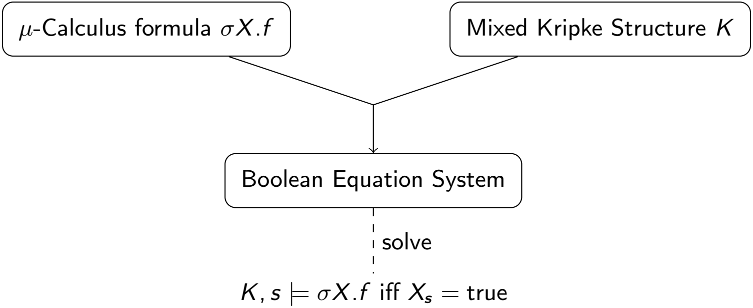
MF 6.073

Outline



- Model checking mu-calculus = solving BES
- Solving BESs conceptually simpler than model checking mu-calculus . still exponential
- BESs are more elementary than mu-calculus . . . . still: fixpoints
- Fixpoints can be understood through an infinite game . . . . Parity games

Parity games

The arena:

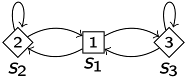
- total graph
- two players: ◇ (Even) and □ (Odd)
- each vertex:
  - has a non-negative priority $p(v)$
  - is owned by one player
- objective: win as many vertices as possible

Definition (Parity game)

A parity game is a four tuple $(V, E, p, (V_\diamond, V_\Box))$ where

- (V, E) is a directed graph
- V a set of vertices partitioned into $V_\diamond$ and $V_\Box$
  - $V_\diamond$: vertices owned by player $\diamond$
  - $V_\Box$: vertices owned by player $\Box$
- E a total edge relation
- $\mathrm{p} : V \to \mathbb{N}$ a priority function

Parity game (example)



$$
\begin{array}{lcl}
V_{\Diamond} & = & \{s_2, s_3\} \\
V_{\Box} & = & \{s_1\} \\
\mathsf{p} & = & \{s_1 \mapsto 1, s_2 \mapsto 2, s_3 \mapsto 3\}
\end{array}
$$

Rules of the game:

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1. place a token on some vertex v
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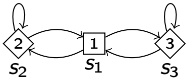
Definition (Winner of a play)

- Let $\pi = v_1 v_2 v_3 \ldots$ be a play
- Let $\inf(\pi)$ be the set of priorities occurring infinitely often in π

Play π is winning for player $\Diamond$ iff $\min(\inf(\pi))$ is even. Likewise for player $\Box$/odd.

Example: winner of a play/winning strategy



Examples of winners of a play:

- Play $(s_1 s_2)^\omega$ won by player $\Box$;
- Play $s_1 s_2^\omega$ won by player $\Diamond$;
- Play $(s_1 s_2 s_1 s_3)^\omega$ won by player $\Box$.

Definition (Strategy)

A strategy for player $\Diamond$ (similarly for $\Box$) is a partial function $\varrho_\Diamond : V^* \times V_\Diamond \to V$

- $v_1 \dots v_n \in V^*$ . . . . sequence of visited vertices (history)
- $v_n \in V_\Diamond$ . . . . vertex owned by $\Diamond$
- $\varrho_\Diamond(v_1 \dots v_{n-1}, v_n) \in \{v \mid (v_n, v) \in E\}$ . . . . rule for moving token from v_n

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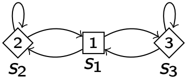
Definition (Consistent plays)

- Let $\pi = v_1 v_2 v_3 \dots$ be an infinite play
- Let $\varrho_\bigcirc$ be a strategy for player $\bigcirc \in \{\Diamond, \Box\}$
- π is consistent with $\varrho_\bigcirc$ iff whenever $\varrho_\bigcirc(v_1 \dots v_{i-1}, v_i)$ is defined, then it is v_{i+1}

$\mathrm{Play}_{\varrho_\bigcirc}(v)$ is the set of all plays starting in v that are consistent with $\varrho_\bigcirc$

Strategy (example)



- possible strategy $\varrho_\Box$: play token from s_1 to s_2 if s_1 has been visited an even number of times, and to s_3 otherwise
- possible strategy $\varrho_\Diamond$ always plays token from s_2 to s_2

Examples of winning strategies:

- $\varrho_\Diamond(\dots, s_2) = s_2$
- $\varrho_\Box(\dots, s_1) = s_3$
- $\varrho_\Diamond(\dots, s_3) = \begin{cases} s_1 & \text{if number of occurrences of } s_3 \text{ is prime} \\ s_3 & \text{otherwise} \end{cases}$

Definition (Winning strategy)

- $\bigcirc \in \{\Diamond, \Box\}$
- $\varrho_{\bigcirc}$ is a strategy for $\bigcirc$

$\varrho_{\bigcirc}$ is a winning strategy from v if every play in $\mathsf{Play}_{\varrho_{\bigcirc}}(v)$ is winning for $\bigcirc$.

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Player $\bigcirc$ wins the vertices in W if from all vertices $v \in W$ she has a winning strategy $\varrho_{\bigcirc}$.

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Natural questions

- Is there always at least one player that can win a vertex?
- Is there a unique winner for each vertex?
- Can the winning strategies be of a particular shape or not?
- Can we compute the winning sets $W_\Diamond$ and $W_\Box$?

Theorem (Positional determinacy)

Player $\bigcirc$ *wins a vertex* w *iff she has a memoryless strategy that is winning from* w

Theorem (Positional determinacy)

Player $\bigcirc$ wins a vertex w iff she has a memoryless strategy that is winning from w

Strategy $\varrho_{\bigcirc}: V^* \times V_{\bigcirc} \to V$ is memoryless (also *history free*) if:

for all histories $\lambda\ v, \lambda'\ v \in V^+$ for which $\varrho_{\bigcirc}$ is defined, we have $\varrho_{\bigcirc}(\lambda, v) = \varrho_{\bigcirc}(\lambda', v)$

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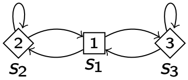
Consequences:

- we can drop the history and consider strategies $\varrho_{\bigcirc}: V_{\bigcirc} \rightarrow V$
- there are only a finite number of memoryless strategies

Memoryless strategy (example)



Let $\varrho_\Diamond(s_2) = s_2$, $\varrho_\Diamond(s_3) = s_1$, and $\varrho_\Box(s_1) = s_3$.

- $\varrho_\Diamond$ is winning from $\{s_2\}$
- $\varrho_\Box$ is winning from $\{s_1, s_3\}$

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Outline

Boolean Equation Systems

Recall Boolean equation systems:

- Boolean expressions: $f, g ::= X \mid \mathsf{true} \mid \mathsf{false} \mid f \wedge g \mid f \vee g$
- Boolean equation system: $\mathcal{E} ::= \varepsilon \mid (\mu X = f)\ \mathcal{E} \mid (\nu X = f)\ \mathcal{E}$

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Lemma ("Tseitin" transformation)

For all Y bound in $\mathcal{E}_0, \mathcal{E}_1$ or $Y = X$:

$$[\mathcal{E}_0\ (\sigma X = f \wedge g)\ \mathcal{E}_1]\eta(Y) = [\mathcal{E}_0\ (\sigma X = f \wedge X')\ (\sigma' X' = g)\ \mathcal{E}_1]\eta(Y)$$

Note: likewise for f, likewise for $f \vee g$

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Lemma (Constant elimination)

For all Y bound in $\mathcal{E}$:

$$[\![\mathcal{E}]\!]\eta(Y) = [\![\mathcal{E}[\mathit{true} := X_{\mathit{true}}]\ (\nu X_{\mathit{true}} = X_{\mathit{true}})]\!]\eta(Y)$$

Note: similarly for false (with $\mu X_{\mathsf{false}} = X_{\mathsf{false}}$)

BES (example)

Consider the following BES:

$$
\begin{array}{lcl}
\mu X & = & X \wedge (Y \vee Z) \\
\nu Y & = & W \vee (X \wedge Y) \\
\mu Z & = & \mathsf{false} \\
\mu W & = & Z \vee (Z \vee W)
\end{array}
$$

This corresponds to the following BES in SRF:

$$
\begin{array}{lcl}
\mu X & = & X \wedge X' \\
\mu X' & = & Y \vee Z \\
\nu Y & = & W \vee Y' \\
\nu Y' & = & X \wedge Y \\
\mu Z & = & X_{\mathsf{false}} \\
\mu W & = & Z \vee (Z \vee W) \\
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Definition (Standard Recursive Form)

A BES is in Standard Recursive Form (SRF) if all right hand sides of Boolean equations adhere to the following syntax:

$$f := X \mid \bigvee F \mid \bigwedge F$$

- X is a proposition variable
- F is a non-empty set of proposition variables

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Observe that:

- all BESs can be transformed into a BES in SRF preserving the solution
- how: repeatedly use "Tseitin" transformation and constant elimination
- the total transformation can be done in polynomial time

Definition (Blocks and ranks)

- a μ-block is a BES of μ-signed equations; likewise: ν-block
- let $\mathcal{E} = \mathcal{B}_1 \cdots \mathcal{B}_n$ for blocks $\mathcal{B}_1, \dots, \mathcal{B}_n$
- Assume for all i, signs of blocks $\mathcal{B}_i$ and $\mathcal{B}_{i+1}$ differ

$$\text{for all } (\sigma X = f) \in \mathcal{B}_i,\ \mathsf{rank}(X) = \begin{cases} i & \text{if } \mathcal{B}_1 \text{ is } \mu\text{-block} \\ i-1 & \text{otherwise} \end{cases}$$

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Observe:

- $\mathsf{rank}(X) = \mathsf{rank}(Y)$ if both X and Y occur in the same block
- $\mathsf{rank}(X)$ is odd iff X is defined in a μ-equation

Rank examples

rank(_)

(1)	μX	$=$	$X \wedge (Y \vee Z)$
(2)	νY	$=$	$W \vee (X \wedge Y)$
(3)	μZ	$=$	false
(3)	μW	$=$	$Z \vee (Z \vee W)$

rank(_)

(1)	μX	$=$	$X \wedge X'$
(1)	$\mu X'$	$=$	$Y \vee Z$
(2)	νY	$=$	$W \vee Y'$
(2)	$\nu Y'$	$=$	$X \wedge Y$
(3)	μZ	$=$	X_{false}
(3)	μW	$=$	$Z \vee (Z \vee W)$
(3)	μX_{false}	$=$	X_{false}

Outline

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Boolean equation systems and Parity games correspond

Let $G = (V, E, p, (V_\Diamond, V_\Box))$ be a parity game

Definition (Parity game to BES)

Define the BES $\mathcal{E}_G$ as follows:

- equations $(\sigma_v X_v = \bigwedge\{X_w \mid (v, w) \in E\})$ for vertices $v \in V_\Box$
- equations $(\sigma_v X_v = \bigvee\{X_w \mid (v, w) \in E\})$ for vertices $v \in V_\Diamond$
- $\sigma_v = \mu$ if $p(v)$ is odd, $\sigma_v = \nu$ otherwise
- ensure $\text{rank}(X_v) \leq \text{rank}(X_u)$ if $p(v) < p(u)$

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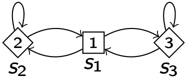
- $\sigma_v = \mu$ if $p(v)$ is odd, $\sigma_v = \nu$ otherwise
- ensure $\text{rank}(X_v) \leq \text{rank}(X_u)$ if $p(v) < p(u)$

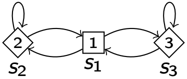
Theorem

Solution to X_v is true $\Leftrightarrow$ player $\diamond$ has winning strategy from v

Parity game to BES example



Parity game to BES example



Corresponds to the following BES:

$$\begin{aligned} \mu X_{s_1} &= X_{s_2} \wedge X_{s_3} \\ \nu X_{s_2} &= X_{s_2} \vee X_{s_1} \\ \mu X_{s_3} &= X_{s_1} \vee X_{s_3} \end{aligned}$$

Parity game to BES intuition

Assume $\mathcal{E}$ is a closed BES in SRF from hereon, unless indicated otherwise.

Lemma

There is a conjunctive BES in SRF $\mathcal{E}'$ constructed from $\mathcal{E}$ by replacing each disjunctive equation $\sigma X_i = \bigvee F_i$ with $\sigma X_i = Y$ for $Y \in F_i$ such that:

$$[\mathcal{E}] = [\mathcal{E}']$$

In the same vein, there is a disjunctive BES in SRF that has the same solution as $\mathcal{E}$.

Parity game to BES intuition

Definition (μ-dominated lasso)

A μ-dominated lasso starting in some X_1 is a finite sequence $X_1 \; X_2 \cdots X_n$, such that:

- We have $X_{i+1} \in F_i$ for $\sigma_i X_i = \bigwedge F_i$ or $\sigma_i X_i = \bigvee F_i$
- We have $X_n \in X_j$ for some $1 \leq j \leq n$.
- $\min\{\mathrm{rank}(X_i) \mid j \leq i \leq n\}$ is odd.

Lemma

Assume $\mathcal{E}$ is conjunctive. Then:

$$[\![\mathcal{E}]\!](X) = \textit{false iff there is a } \mu\textit{-dominated lasso starting in } X$$

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Parity game to BES intuition

Theorem

Solution to X_v is true $\Leftrightarrow$ player $\Diamond$ has winning strategy from v

Proof.

- $\Leftarrow$
  - Assume player $\Diamond$ has a winning strategy ϱ from vertex v.
  - Let $\mathcal{E}$ be the BES obtained from the parity game.
  - Construct $\mathcal{E}'$ from $\mathcal{E}$ by replacing every disjunctive equation as follows:

$$(\sigma X_u = \bigvee F) \text{ becomes } (\sigma X_u = X_{\varrho(u)})$$

  - Towards a contradiction, suppose $[E'](X_v) = \text{false}$
  - Then there must be a μ-dominated lasso starting in X_v
  - But that means that the lowest rank on the lasso is odd
  - Hence, by the transformation, there must be an infinite path in the parity game on which the lowest priority is odd
  - Hence, ϱ is not winning for $\Diamond$. Contradiction
- $\Rightarrow$
  - Dually, assume $\Box$ has a winning strategy and prove $[\mathcal{E}](X_v) = \text{false}$.

$\Box$

Let $\mathcal{E}$ be a closed BES in SRF.

Definition (BES to parity game)

Define a parity game $G_{\mathcal{E}} = (V, E, p, (V_{\Diamond}, V_{\Box}))$ as follows:

- $v_X \in V$ iff there is an equation for X in $\mathcal{E}$
- $(v_X, v_Y) \in E$ iff propositional variable Y occurs in f in $\sigma X = f$
- $p(v_X) = \text{rank}(X)$ for all equations $(\sigma X = f)$ in $\mathcal{E}$
- $v_X \in V_{\Box}$ iff the equation for X is of the form $(\sigma X = \bigwedge F)$
- $V_{\Diamond} = V \setminus V_{\Box}$

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Boolean equation systems and Parity games correspond

Let $\mathcal{E}$ be a closed BES in SRF.

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- $v_X \in V_\Box$ iff the equation for X is of the form $(\sigma X = \bigwedge F)$
- $V_\diamond = V \setminus V_\Box$

Theorem

Player $\diamond$ has winning strategy from v_X $\Leftrightarrow$ the solution of X is true

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BES vs parity game (example)

Consider the following BES:

$$
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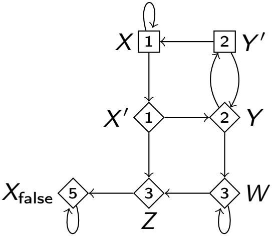
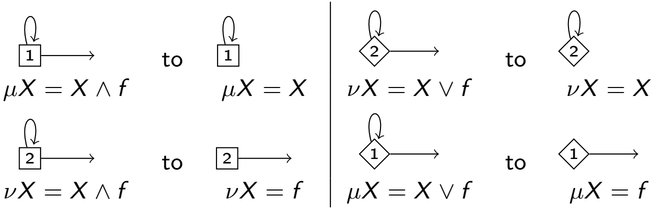
Its parity game is:

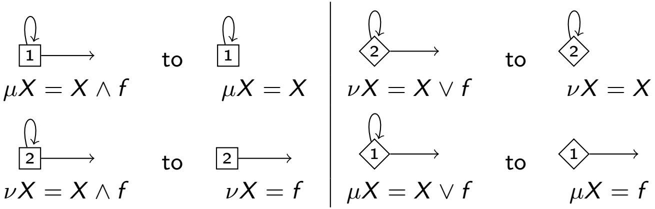


Outline

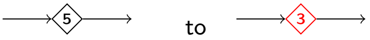
Self-loop elimination



Self-loop elimination



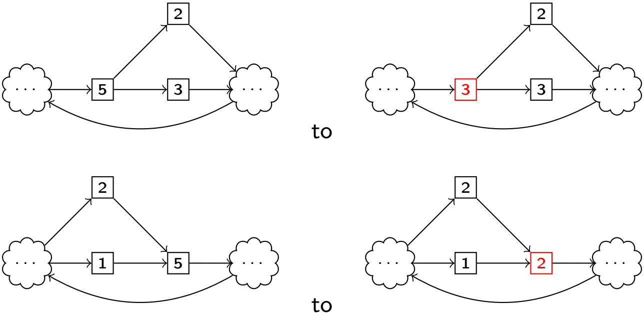
Priority compaction



In case priority 4 does not occur in the parity game. Evenness must be preserved!

Simplifying parity games

Priority propagation



Corresponds to re-ordering of equations in BES, which is generally unsafe!

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Outline

Summary

- Computing winners in parity games = solving BESs
- Reduction parity games $\leftrightarrow$ BESs is polynomial
- Operational interpretation of fixpoints:
  - μ-fixpoint: odd priorities; can only be won by $\Diamond$ if it ensures stretches are finite
  - ν-fixpoint: even priorities; benign for player $\Diamond$
- Simplifications
- No algorithm yet . . . . but

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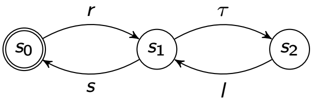
Next week:

- Recursive algorithm

Outline

Exercise



Consider the following modal μ-calculus formula f:

$$\nu X.([r]X \wedge ((\nu Y.\langle\tau\rangle Y \vee \langle l\rangle Y) \vee (\mu Z.(([l]Z \wedge [s]Z) \vee \langle s\rangle \mathrm{true}))))$$

- Translate the model checking question $M \models f$ to a BES.
- Transform the resulting BES into a parity game.
- Determine whether f holds in s_0 by solving the obtained parity game, and
- provide a winning strategy that justifies this solution.