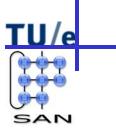
Reinder J. Bril, r.j.bril@tue.nl TU/e Informatica, System Architecture and Networking

1



Module Architecture Control using Relation Algebra

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18-11-2009

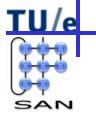
Questions

- Development artefacts of a system?
- How to maintain consistency:
 - between these artefacts (inter)?
 - between elements of an artefact (intra)?
- This lecture:
 - consistency between architecture and implementation
- SAN:

[]/e

- consistency of designs [Lange et al 05]
- generalization [Muskens et al 05].

Goals



Student understands:

- the need for module architecture control;
- how module architecture control can be performed;
- the basics of relation algebra.
- Student can apply relation algebra on a simple example.

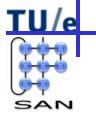


Overview

- TU/e
 - Context and motivation
 - Software architecture recap
 - Application domain
 - Module architecture notions
 - Relation algebra
 - Verification
 - Conclusion
 - References



Overview



Context and motivation

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#

Software architecture – recap

TU/e

Stakeholders

- End-User:
 - behavior, performance, security, reliability
- Customers:
 - low cost, timely delivery
- **Product-Management**:
 - features, short time to market, low cost, parity with products

Development:

- low cost, employability
- Maintenance:
 - modifiability

SAN

Software architecture – recap

End-User:

behavior, performance, security, reliability

Customers:

□ low cost, timely delivery

Product-Management:

features, short time to market, low cost, parity with products

• Development:

□ low cost, employability

• Maintenance:

modifiability

roduct view

Stakeholders

TU/e

H

Software architecture – recap

- End-User:
 - behavior, performance, security, reliability
 - Customers: □ low cost, timely delivery
 - **Product-Management**:
 - features, short time to market, low cost, parity with products

Development:

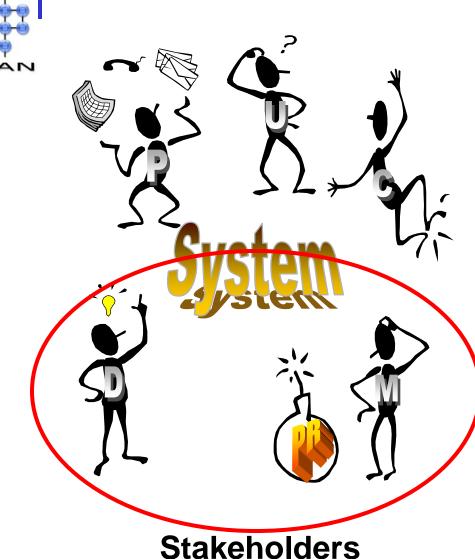
□ low cost, employability

Maintenance:

modifiability

evelopment viev

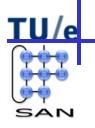
8



TU/e



Software architecture – recap



Software architecture:

- earliest artifact \square
- means for mutual communication
- enables analysis of concerns \square
- manifests concerns as system qualities
- Software architecture is vital ! ٠

[Bass et al 1995]

Software architecture – recap

'll/e

٠

Need for (system) architecture

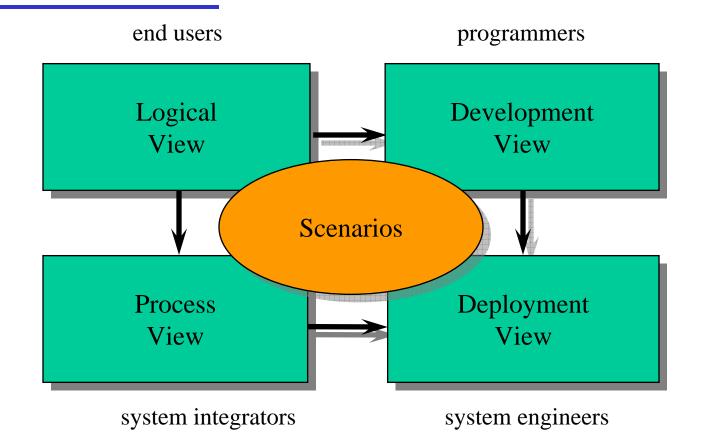
If a project has not achieved a system architecture, including its rationale, the project should not proceed to full-scale system development. Specifying the architecture as a deliverable enables its use throughout the development and maintenance process.

[Boehm 1995]



TU/e

Software architecture – recap

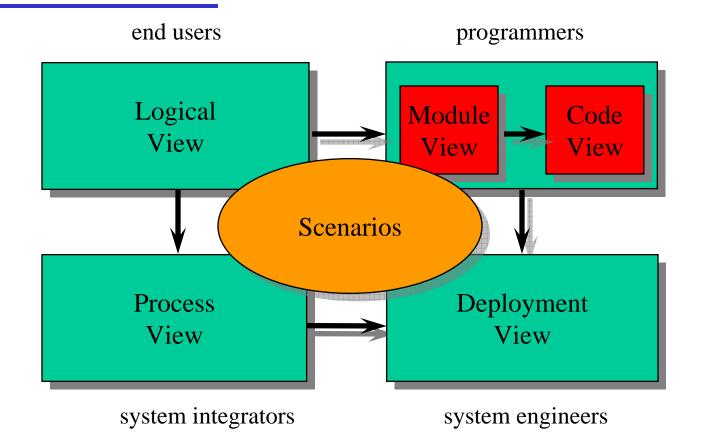


4+1 View Model [Kruchten 95]



TU/e

Software architecture – recap

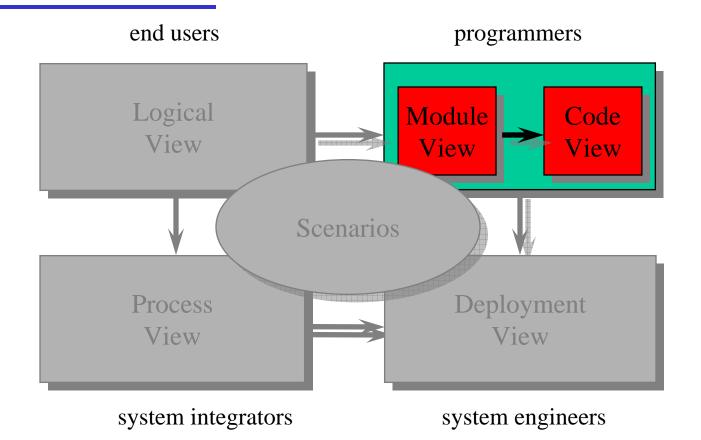


4+1 View Model Revisited



TU/e

Software architecture – recap



4+1 View Model Revisited



Software architecture – recap TU/e programmers Code Module View View

Development View



Software architecture – recap

- A software architecture characterization:
 - Components (or [architectural] entities);
 - Connections;
 - Constraints.
- Module view: ۲

TII/e

- System, Subsystems, Components;
- Part-of relation and uses relation;
- Layering, orthogonally.
- Code view:
 - Directories, Files;
 - Directory structure, location of files, and include relation; ____
 - File name conventions and file length constraints.



Overview

- Context and motivation
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TU/e

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- Telecommunications domain
- SOPHO: Philips' family of PBXs
 - 100 1M telephony lines

ΓU/e

- origin dating back to early 1980's
- economic lifetime ~ 15 years
- maintenance obligations \geq 10 years
- 5 K files, 2.5 MLOC in C++
- successful \Rightarrow asset
 - \Rightarrow careful to maintain this legacy

- Complexity of legacy systems
 - hard to understand (e.g. size);
 - documentation out-of-date;
 - gap between intrinsic and experienced complexity.
- Architecture vital, but not maintained ...:
 - recovery;

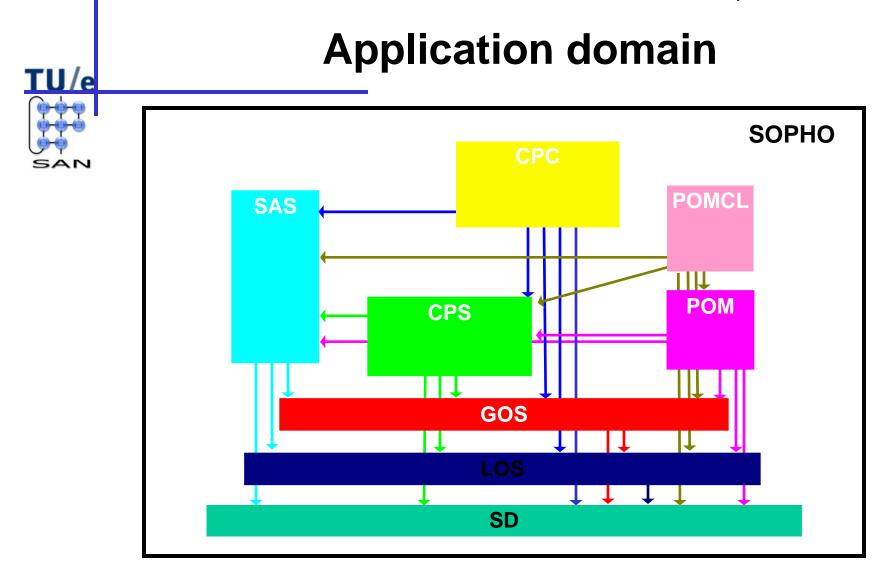
TU/e

- analysis;
- verification and control;
- (improvements).
- Need for architectural support !

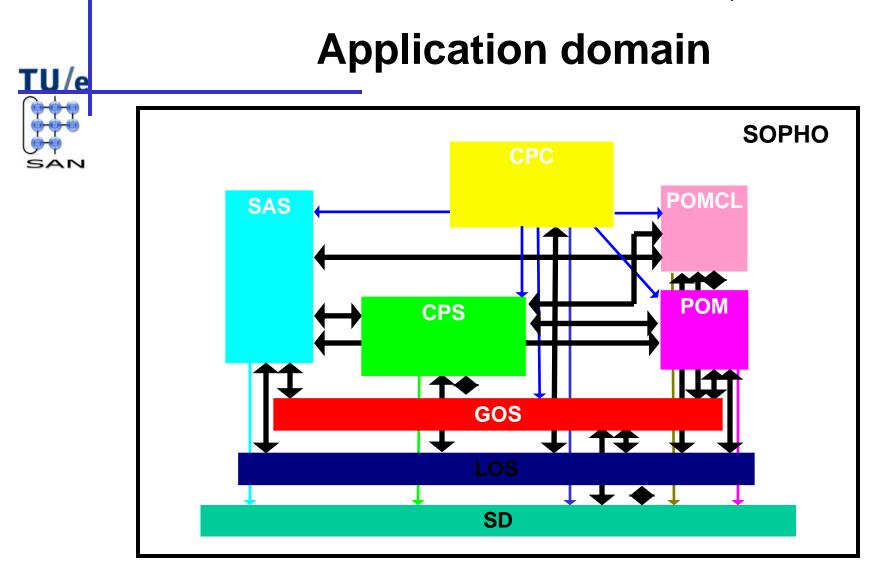
- Characteristics (development view):
 - module view

[]/e

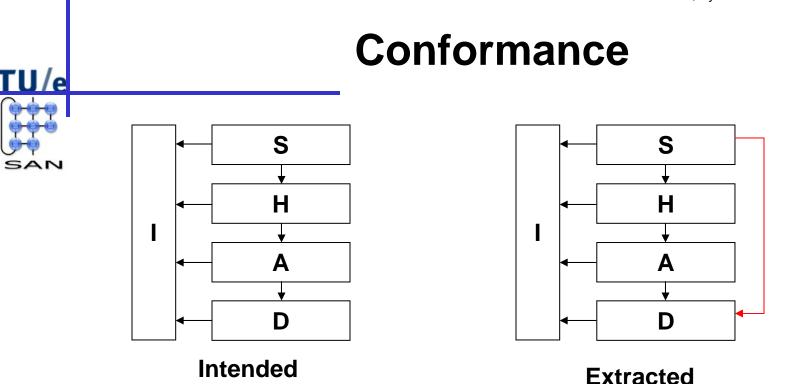
- 8 K architectural entities;
- organised in an unbalanced tree, depth 5 12;
- layered system, consisting of 8 subsystems.
- code view
 - 1 directory with 5 K files, and 2.5 MLOC in C++;
 - 35 K include statements;
 - File names based on "12 NCs", file length varies from 100 to 20 K lines.



"Intended" module architecture (documentation + software architects)



"Derived" module architecture (extracted from the implementation)



Causes when "intended" and "extracted" differ:

- 1. "intended" is wrong (e.g. out-of-date): improve;
- 2. "extracted" is wrong: improve;
- 3. implementation is optimized for, e.g., speed \Rightarrow refinement.

- Ensure conformance to an architecture !
 - Keep the architecture up-to-date

[]/e

- Approach using relation algebra (RPA):
 - Represent the "intended" architecture in RPA.
 - Extract the "derived" architecture from the implementation, and represent in RPA.
 - Express "conformance" in RPA.
 - Ensure conformance by means of verification (using RPA) and improvements (i.e. control).

Overview

- Context and motivation
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TU/e

SAN

- Conclusion
- References



Overview

- Context and motivation
- Module architecture notions
 - Module diagram
 - Decomposition tree
 - Lifting

TU/e

- Hiding
- Lowering
- Weights
- Relation algebra
- Verification
- Conclusion
- References

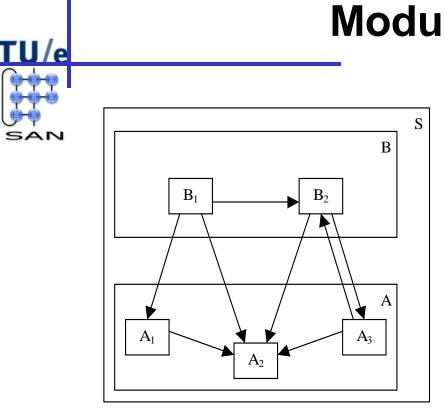


Module architecture notions

TU/e SAN

Module diagram

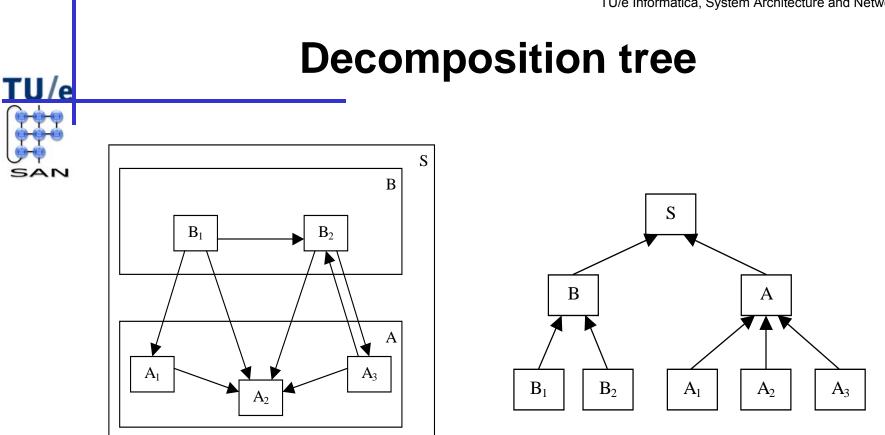
- **Decomposition tree**
- Lifting
- Hiding
- Lowering
- Weights



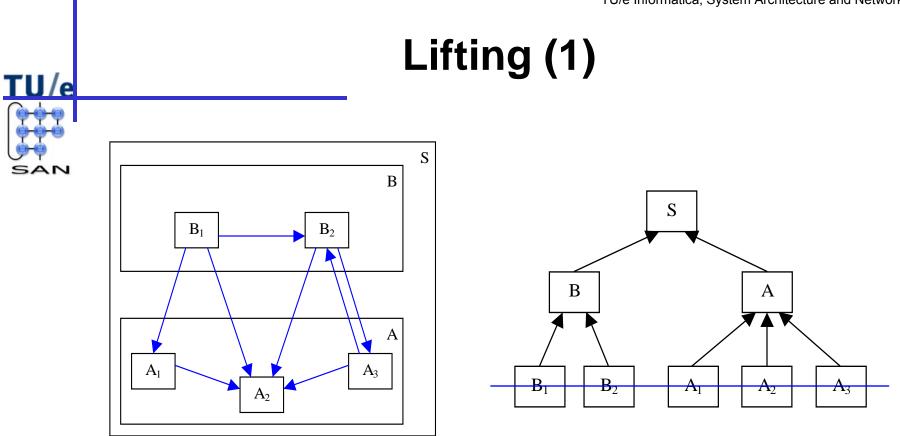
Module diagram

- Visualises a system's architecture
- "Boxes-in-boxes" representation
- Boxes represent entities
- Arrows represent dependencies

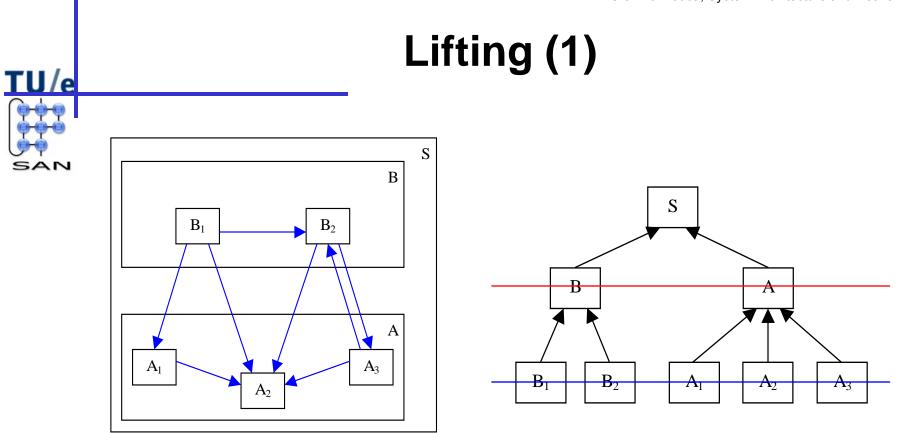
System S is not layered



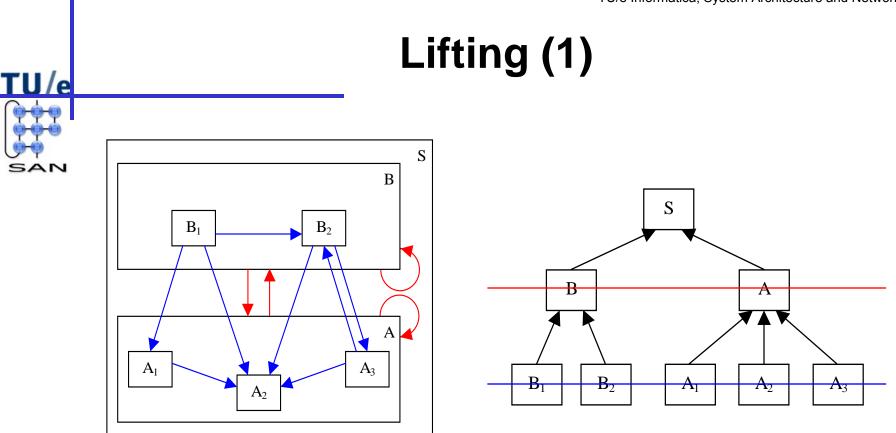
System S is *balanced, and* the decomposition tree has *3 levels*



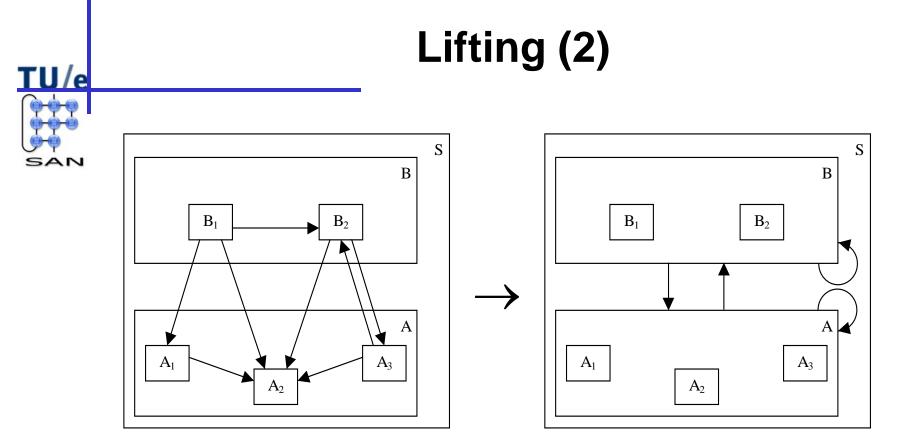
Uses relation corresponds with a level (*tree-cut*)

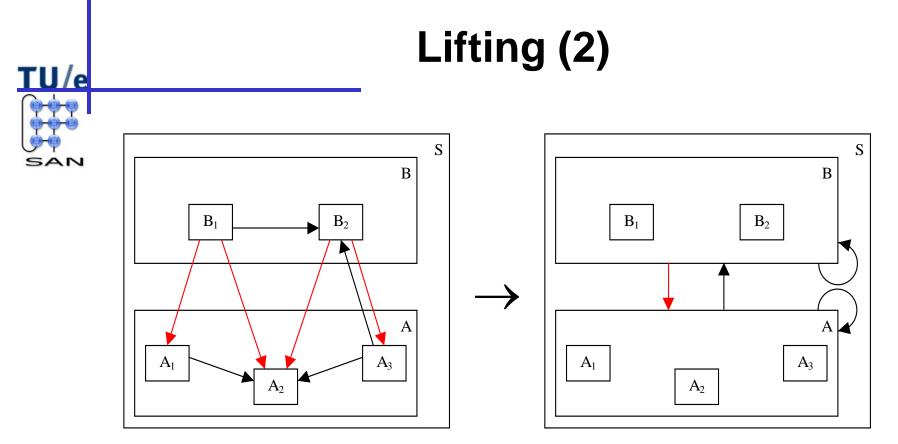


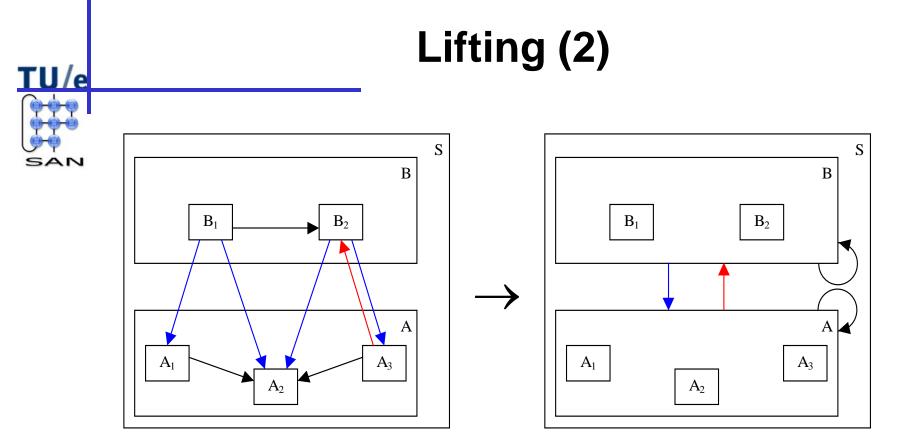
Transform a relation to a higher level.

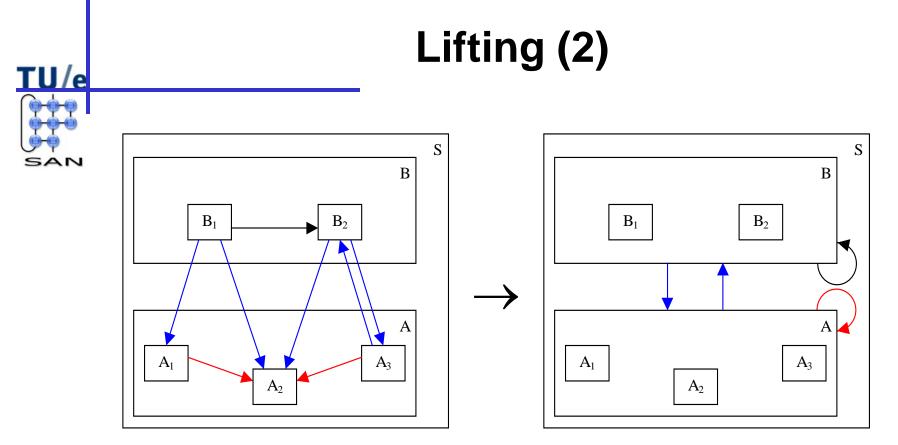


Transform a relation to a higher level, i.e. *replace both the source and the destination of each arrow by its enclosing entity.*

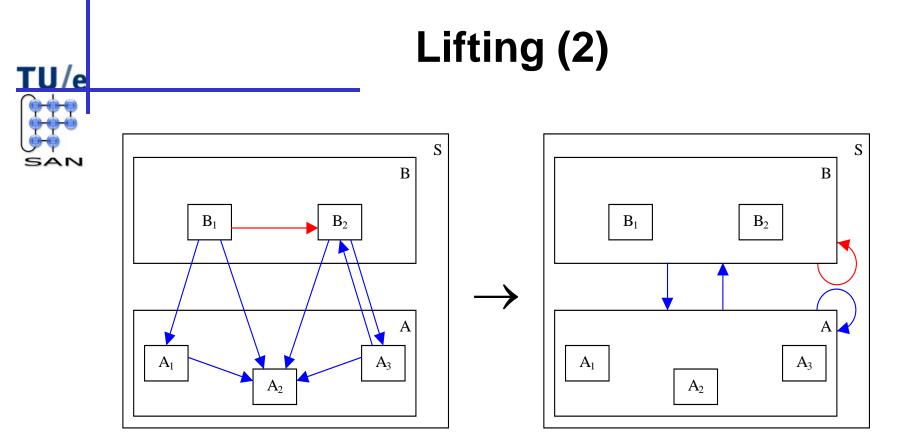


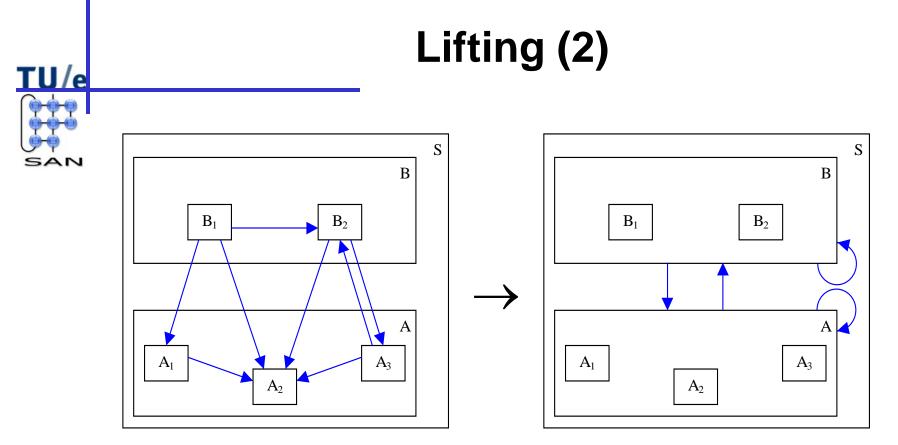


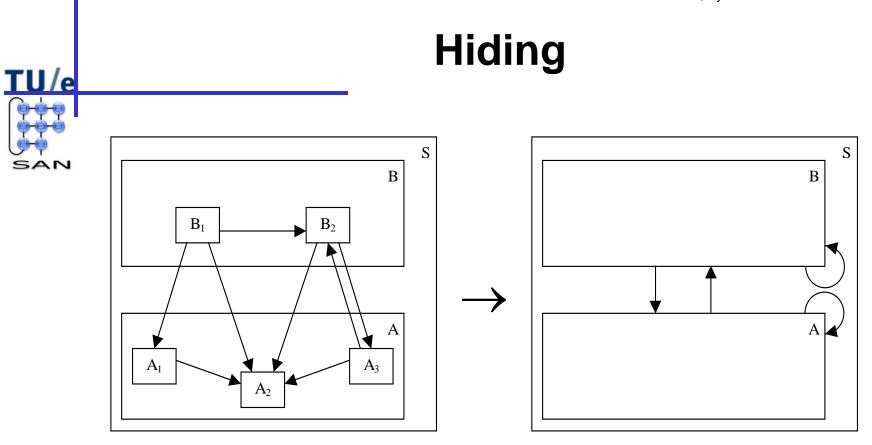




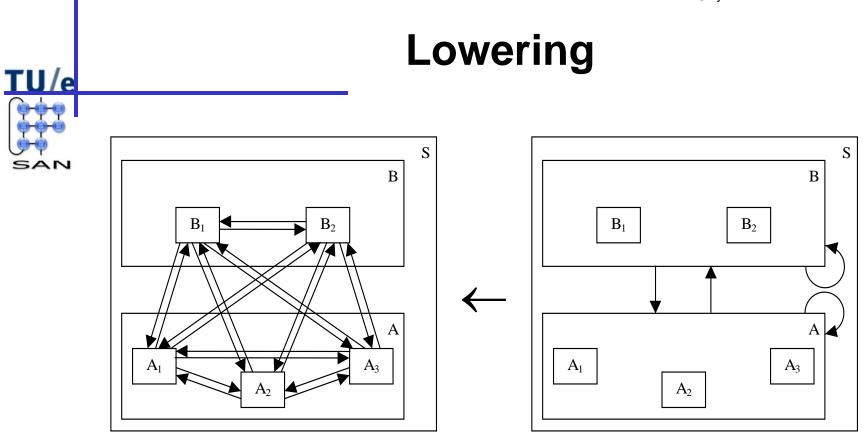
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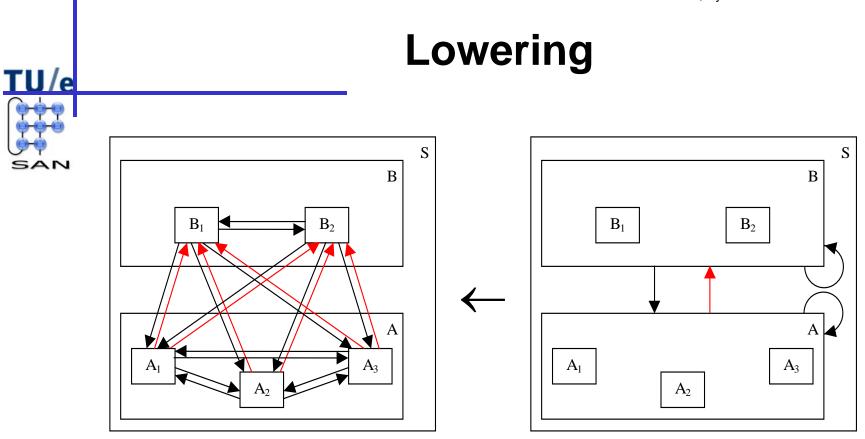




Hiding the decomposition structure of both A and B

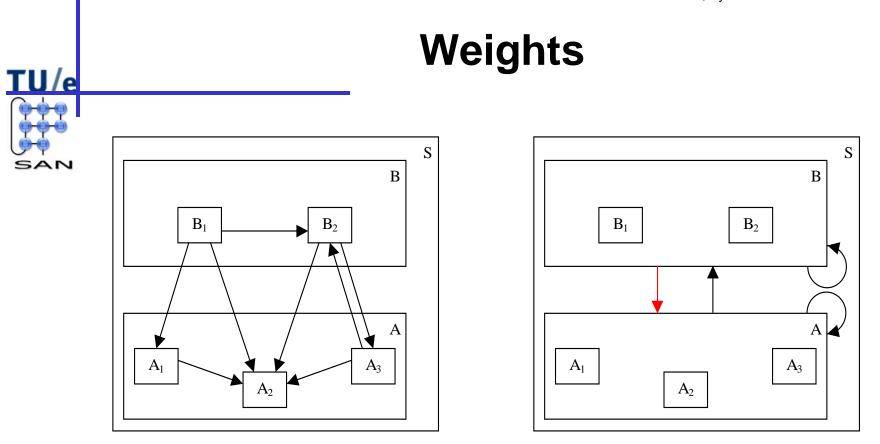


both cases a complete graph.

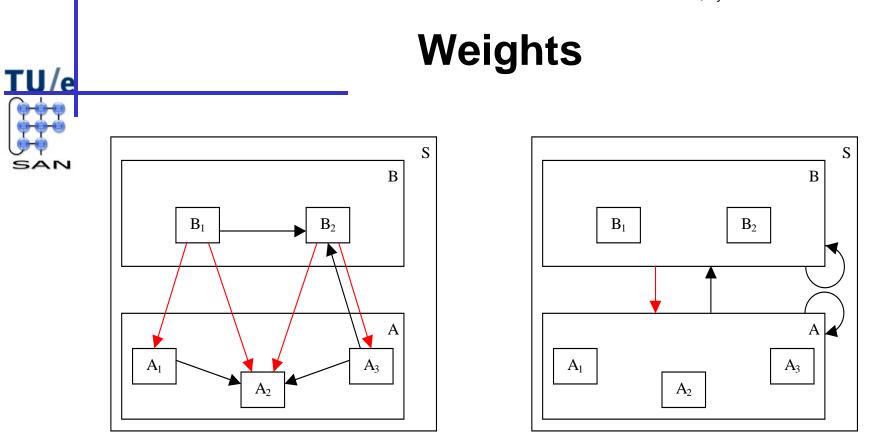


Application: architectural verification,

e.g. layering

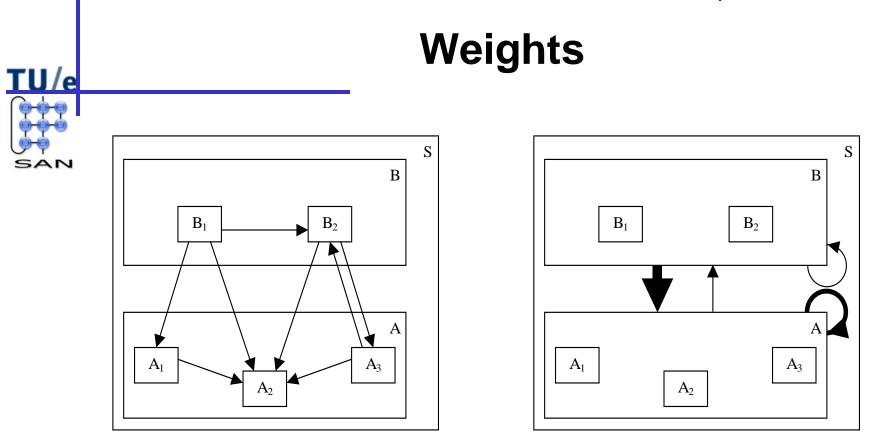


Which value to be associated ?



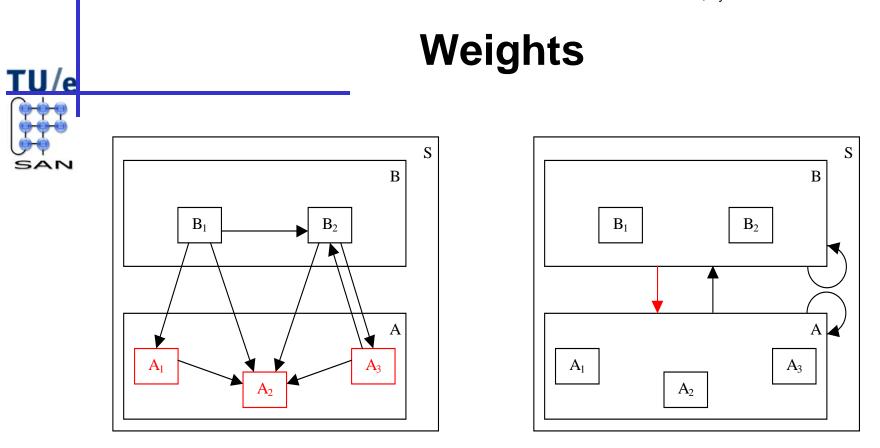
4: number of uses relations

size-oriented weight

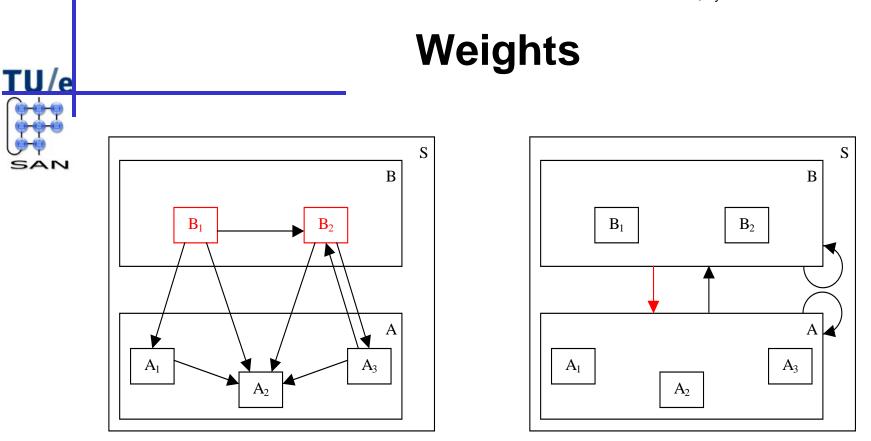


Fisheye view of the

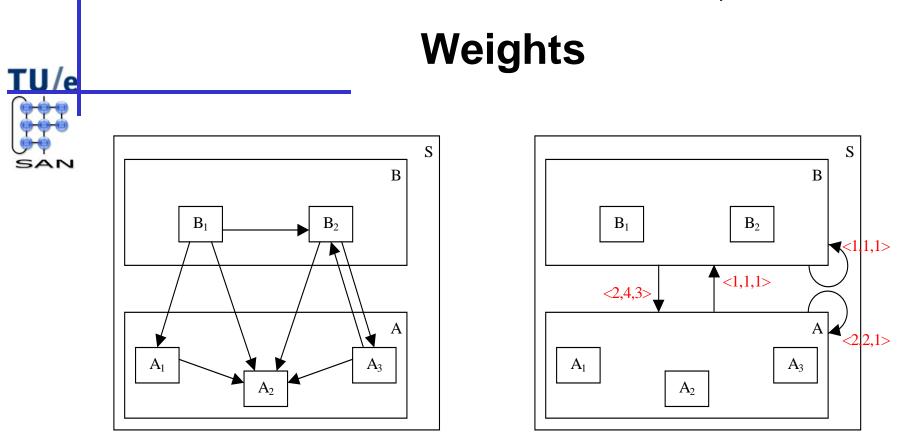
size-oriented weight



3: number of used entities (A₁, A₂ and A₃) *fan-in-oriented weight*



2: number of *using entities* (B₁ and B₂) fan-out-oriented weight



Each weight has its merits during architectural analysis



Overview

- Context and motivation
- Module architecture notions •
- **Relation algebra** •
 - Usage

TU/e

SAN

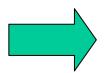
- Overview
- Examples
- Application
- Verification
- Conclusion
- References

Relation Algebra: Usage

- Visualisation and view calculations
 - reverse architecting purposes
- Relational calculus

[]/e

- software architecture analysis
- Architectural rules
 - software architecture verification
- Architectural metrics
 - software architectural quality assurance
- (formal basis of tools)

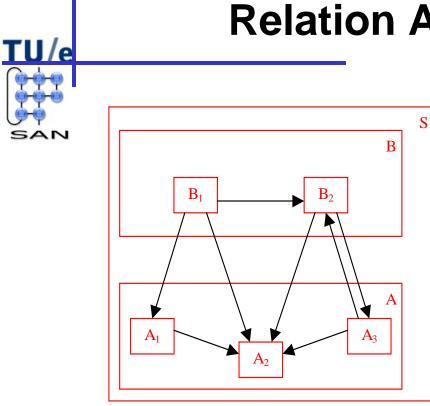






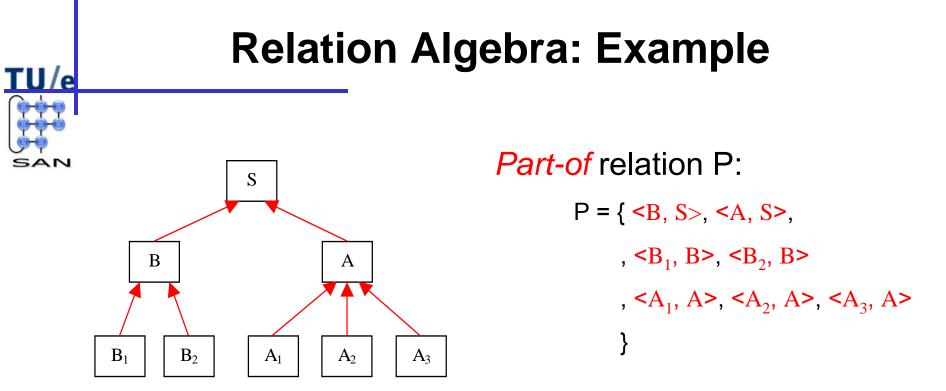
TU/e

- $\emptyset, \cup, \cap, -, \dots$
- Relations: sets of pairs
 - -; (composition), -1, +, *, $|,\uparrow$ (lifting), \downarrow (lowering)
- Multi-sets: bags
 [], []
- Multi-relations: bags of pairs



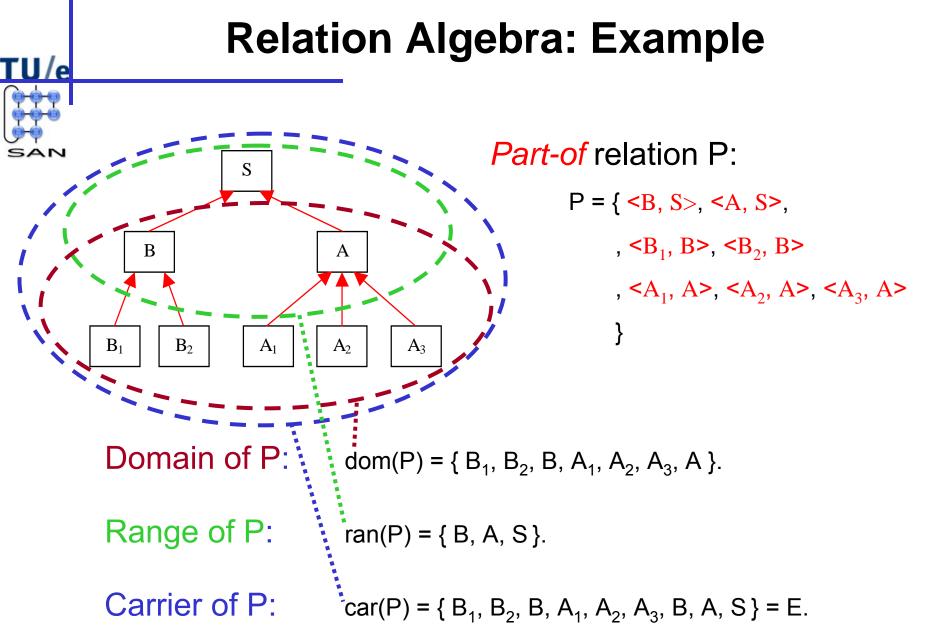
Relation Algebra: Example

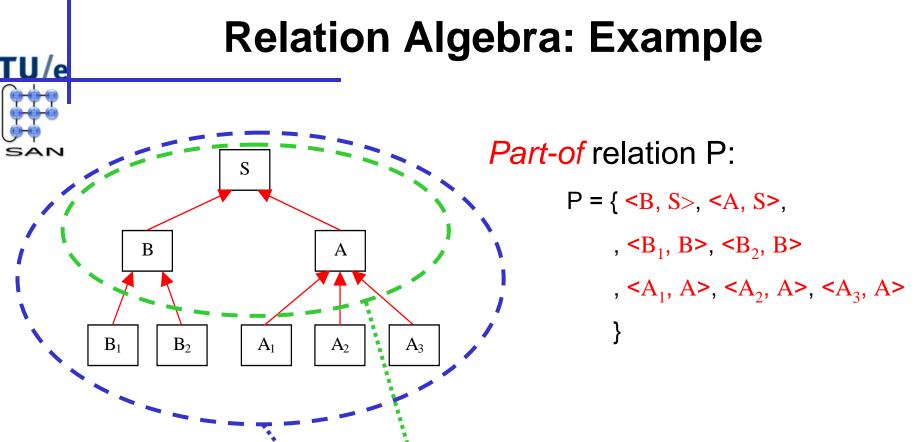
Set of *Entities* E: E = { S, A, A₁, A₂, A₃, B, B₁, B₂}



A part-of relation:

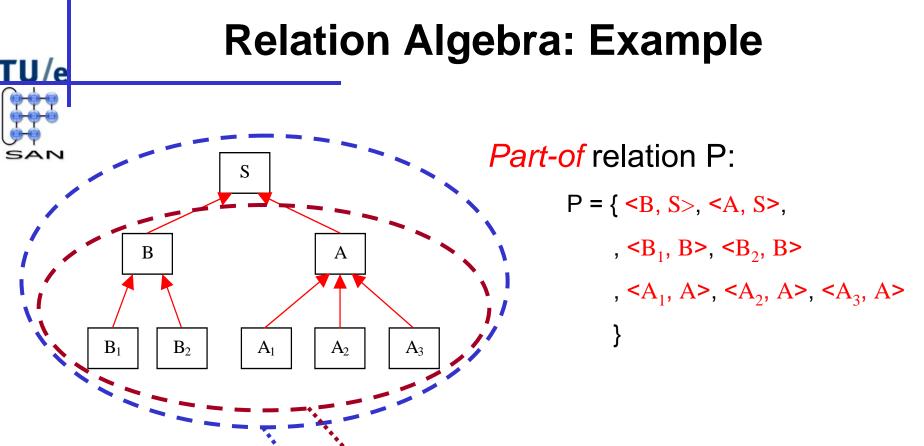
- describes the decomposition tree;
- is both: *functional* and *a-cyclic*.





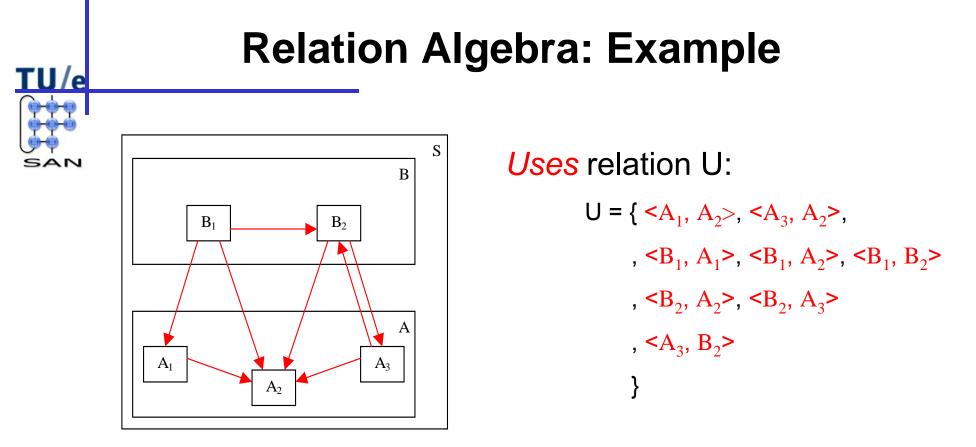
Question: How to express the leafs of the decomposition tree in relation algebra using P?

Answer: leafs(E) = car(P) - ran(P) = { B_1, B_2, A_1, A_2, A_3 }.



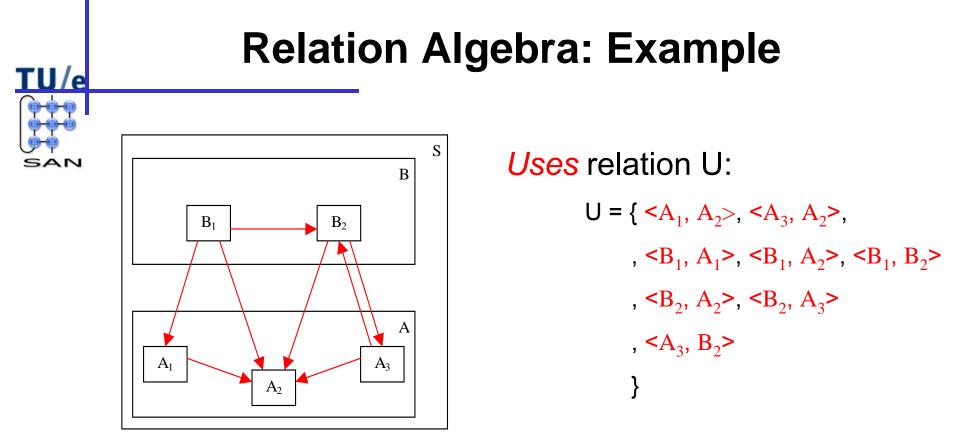
Question: How to express the root of the decomposition tree in relation algebra using P?

Answer: $root(E) = car(P) - dom(P) = \{S\}.$



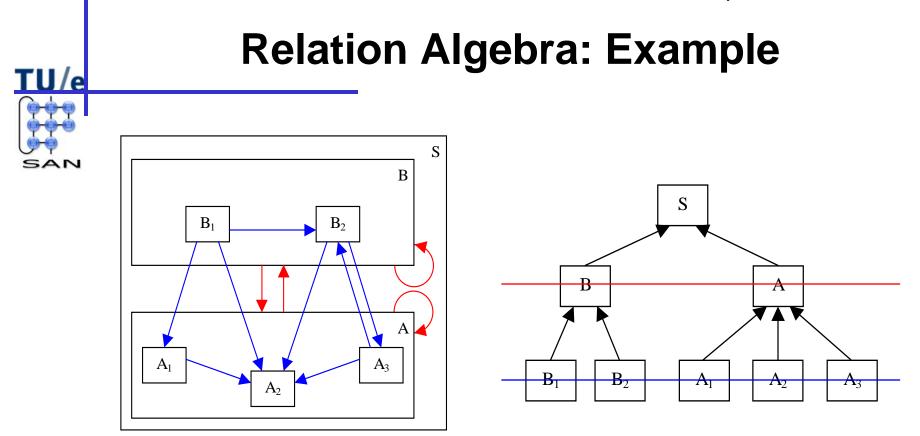
Question: How to express the entities that use, but are not used by, entities in relation algebra ?

Answer: use_only = car(U) - ran(U) = $\{B_1\}$.

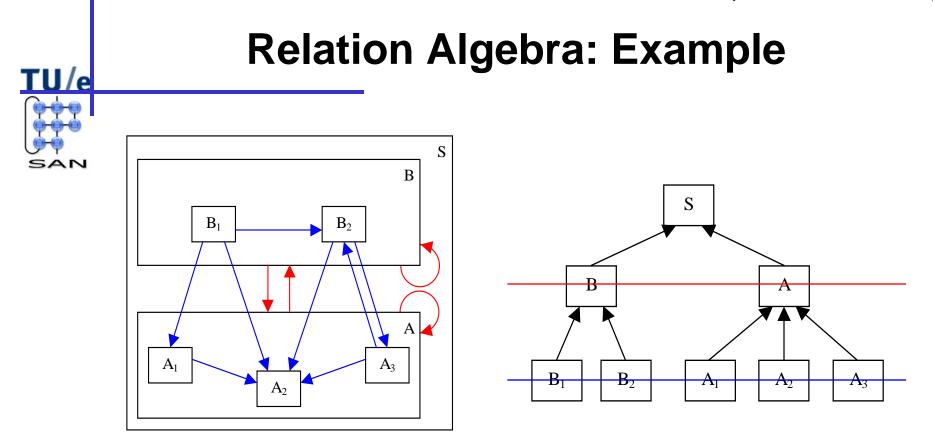


Question: How to express the leaf entities that neither use nor are used by entities in relation algebra ?

Answer: isolated = leafs(E) – car(U) = \emptyset .



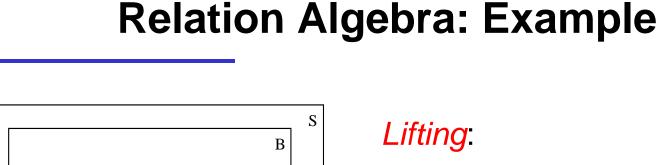
Transform a relation to a higher level, i.e. *replace both the source and the destination of each arrow by its enclosing entity*.



Example: $\langle B_1, A_1 \rangle$

- *replace source:* B₁ by B;
- replace destination: A_1 by A.





А

 A_3

 B_2

 A_2

 B_1

 A_1

TU/e

def.: $U \uparrow P \equiv P^{-1}$; U; P

$$ex.: \quad {}^{\mathsf{C}}\mathsf{B}_1, \mathsf{B}^{\mathsf{C}} \in \mathsf{P} \Leftrightarrow {}^{\mathsf{C}}\mathsf{B}, \mathsf{B}_1^{\mathsf{C}} \in \mathsf{P}^{-1};$$

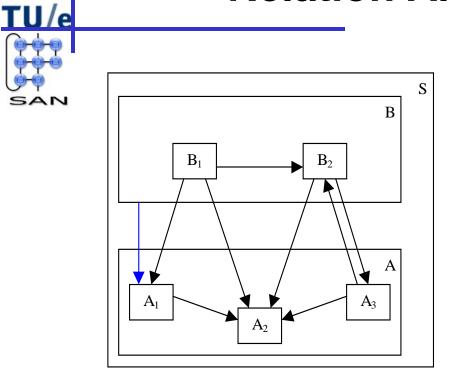
$$< B_1, A_1 > \in U;$$

$$A_1, A \in P;$$

For $U \uparrow P$, P must be a part-of relation







Lifting.

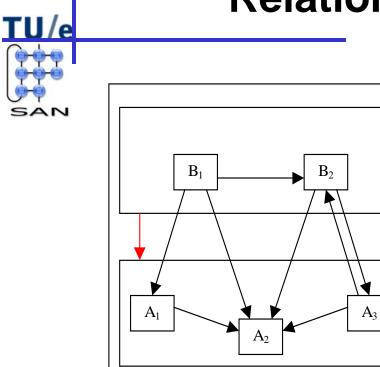
def.: $U \uparrow P \equiv P^{-1}$; U; P

$$ex.: \quad \mathsf{<}B_1, B\mathsf{>} \in \mathsf{P} \Leftrightarrow \mathsf{<}B, B_1\mathsf{>} \in \mathsf{P}^{-1};$$

 $< B_1, A_1 > \in U;$

$$A_1, A \in \mathsf{P};$$

<B, A_1 **>** \in P⁻¹ ; U



Relation Algebra: Example

S

В

А

Lifting.

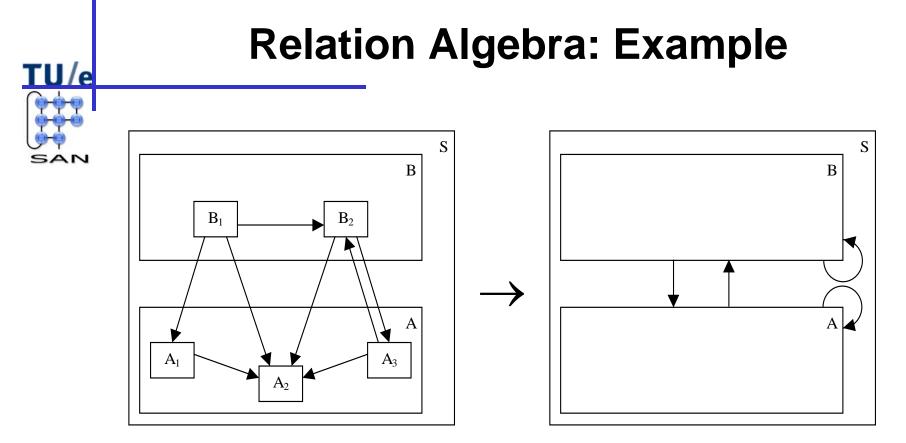
def.: $U \uparrow P \equiv P^{-1}; U; P$

ex.:
$$\langle B_1, B \rangle \in P \Leftrightarrow \langle B, B_1 \rangle \in P^{-1};$$

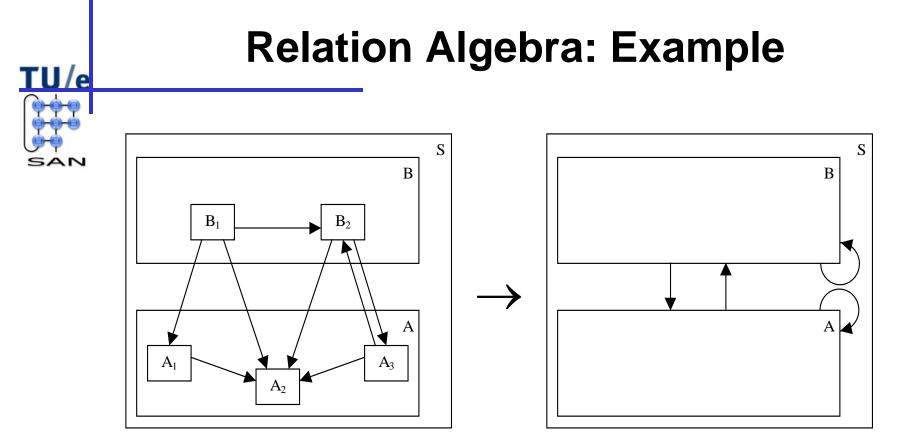
 $\mathsf{B}_1, \mathsf{A}_1 \mathsf{E} \mathsf{U};$

$$A_1, A \in P;$$

<B, A> ∈ P⁻¹; U; P

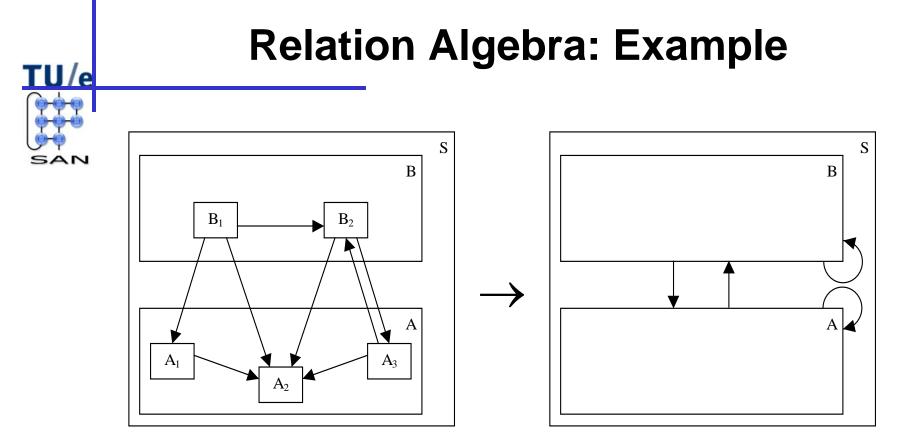


Hiding the decomposition structure of both A and B



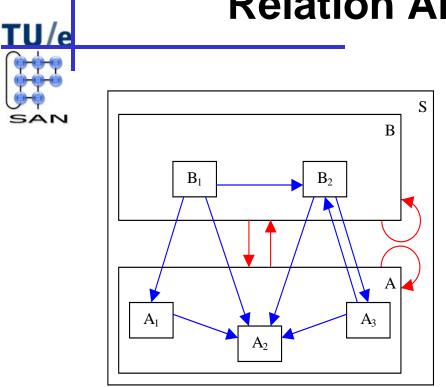
Question: How to express the set of entities E' after hiding in relation algebra using P?

Answer: E' = ran(P) = { S, A, B }.



Question: How to express the part-of relation P' after hiding in relation algebra ?

Answer: $P' = P |_{car} E' = \{ S, A, B \}.$

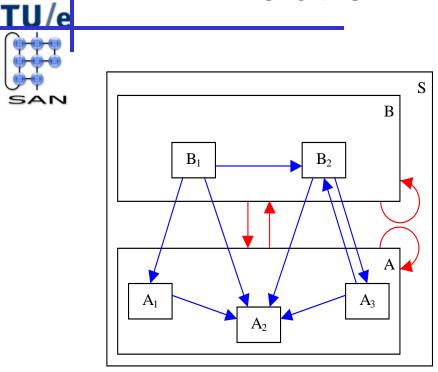


Relation Algebra: Example

Lowering:

def.: $U \downarrow P \equiv P$; U; P⁻¹

For U \downarrow P, P must be a part-of relation



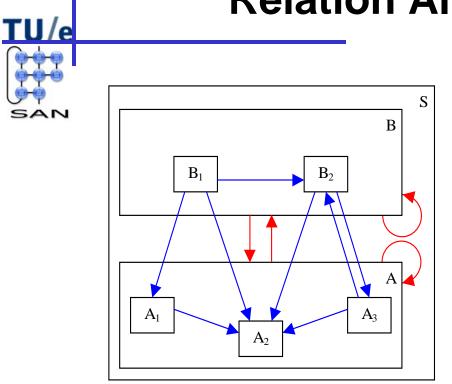
Relation Algebra: Example

Lowering: def.: $U \downarrow P \equiv P ; U ; P^{-1}$ ex.: $\{<A, B>\} \downarrow P = \{<A_1, B_1>, <A_1, B_2>$ $,<A_2, B_1>, <A_2, B_2>$ $,<A_3, B_1>, <A_2, B_2>$ $\}$ $\{<A, B>\} \downarrow P \cap U = \{<A_3, B_2>\}$

Layering rule:

 $\{\langle A, B \rangle\} \downarrow P \cap U = \emptyset$

SAN



Relation Algebra: Example

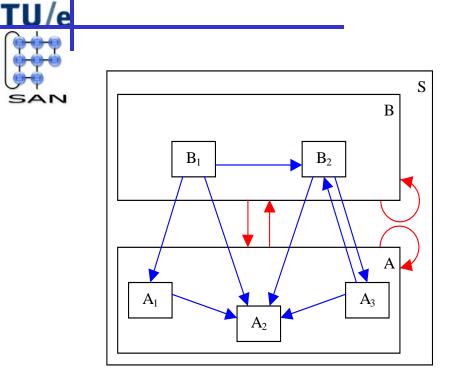
Size-oriented lifting.

def.: $U_{M} \uparrow P \equiv \lceil P^{-1} \rceil; U_{M}; \lceil P \rceil$ where $U_{M} = [U]$

ex.:
$$\langle B, A \rangle$$
 4 times in U_M \uparrow P

₿ ‡ SAN

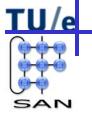




Fan-out-oriented lifting. $U_{M} \uparrow \triangleright P \equiv [P^{-1}]; [U_{M}]; P]$ <B, A> 2 times in $U_{M} \uparrow \triangleright P$ Fan-in-oriented lifting. $U_{M} \uparrow \lhd P \equiv \left\lceil P^{-1}; \left\lfloor U_{M} \right\rfloor \right\rceil; \left\lceil P \right\rceil$ $\langle B, A \rangle$ 3 times in U_M $\uparrow \triangleleft P$

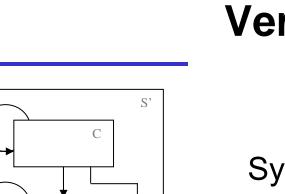


Relation Algebra: Application



- Analytical techniques:
 - lifting,
 - checking of rules,
 - impact analysis,
 - finding unused and unavailable components,
 - identification of top and bottom layers,
 - study of alternative component groupings,
 - improvement of presentation by suppressing elements,
 - improvement of locality,
 - checking of linearity,
 - analysis of cycles,
 - improvement of presentation by transitive reduction.

[Feijs et al 98]



В

Α

TU/e

Verification

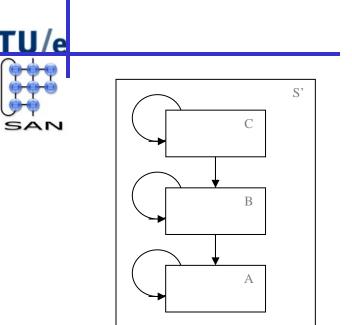
System S' consists of 3 levels, and is *layered*.

Question: How to express layering rule LR for S' in relation algebra?

Answer: Layering rule LR for S':

LR(S'): {<A, B>, <A, C>, <B, C>} \downarrow P \cap U = ø



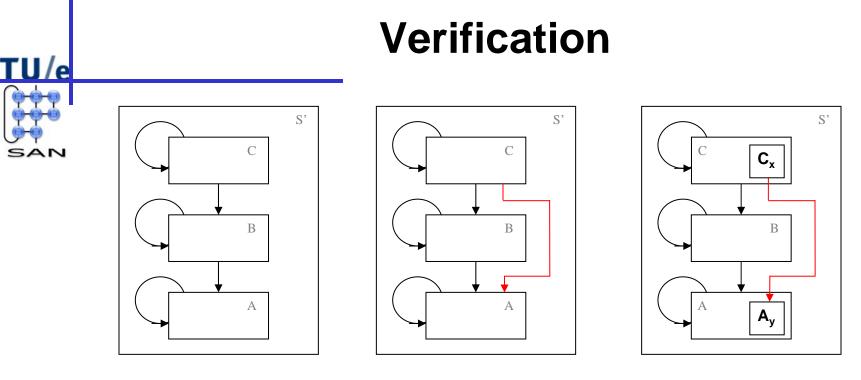


System S' consists of 3 levels, and is *strictly* layered.

Question: How to express strictly layering rule SLR for S' in relation algebra?

Answer: Strictly layering rule SLR for S':

SLR(S'): {<A, B>, <A, C>, <B, C>, <C, A>} ↓ P ∩ U = ∅



Intended

Derived

Question: How to express the exception to the strictly layering rule SLR(S') for S' in relation algebra?

Answer: Delete specific uses relation between modules !

 $\{<A, B>, <A, C>, <B, C>, <C, A>\} \downarrow P \cap U - \{<C_x, A_y>\} = \emptyset$

Conclusion

Goals revisited

TU/e

- Student understands:
 - the need for module architecture control;
 - diversion of intended and derived architecture over time
 - how module architecture control can be performed;
 - describe the architecture in terms of rules, check compliance of the derived architecture, and adapt
 - the basics of relation algebra.
- Student can apply relation algebra on a simple example.

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TU/e

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