

“Cheat sheet” Exams 2WC09/2WC12, Fall 2012

At the exams, use of a simple, non-programmable pocket calculator is allowed. All other electronic equipment is not allowed, nor any notes or books.

The following formulas will be provided at the exams.

Shannon entropy of r.v. X taking values in $\{x_1, \dots, x_n\}$ with probability $p_i = \Pr[X = x_i]$:

$$H(X) = - \sum_{i=1}^n p_i \log_2 p_i$$

Digital Signature Algorithm (DSA), for hash function h .

Key generation. Given primes p, q such that $p|(q-1)$ and an element g of order p in \mathbb{Z}_q^* , choose private key m_U uniform at random in \mathbb{Z}_p^* , and set public key $c_U = g^{m_U} \bmod q$.

Signature generation. Given message M , choose r uniform at random in \mathbb{Z}_p^* , and set

$$R = (g^r \bmod q) \bmod p, \quad S = (h(M) + m_U R)r^{-1} \bmod p$$

Signature verification. Given message M , signature (R, S) , check if

$$R = (g^x c_U^y \bmod q) \bmod p, \quad \text{where } x = h(M)S^{-1} \bmod p, y = RS^{-1} \bmod p$$

Addition formulas for elliptic curve E over any finite field of characteristic p .

Let $P_3 = P_1 + P_2$.

If $P_1 = \mathcal{O}$, then $P_3 = P_2$.

If $P_2 = \mathcal{O}$, then $P_3 = P_1$.

If $P_1 = -P_2$, then $P_3 = \mathcal{O}$.

Otherwise, let $P_1 = (x_1, y_1)$, $P_2 = (x_2, y_2)$, then $P_3 = (x_3, y_3)$ is given as follows.

$$\begin{array}{l|l} \text{Case } p = 2 & \text{Case } p \neq 2 \\ E : y^2 + xy = x^3 + ax^2 + c & E : y^2 = x^3 + ax^2 + bx + c \\ \begin{array}{l} x_3 = x_1^2 + \frac{c}{x_1^2} \\ y_3 = x_1^2 + (x_1 + \frac{y_1}{x_1})x_3 \end{array} & \begin{array}{l} x_3 = \lambda^2 - a - x_1 - x_2 \\ y_3 = \lambda(x_1 - x_3) - y_1 \end{array} \end{array} \quad \lambda = \begin{cases} \frac{3x_1^2 + 2ax_1 + b}{2y_1}, & x_1 = x_2 \\ \frac{y_2 - y_1}{x_2 - x_1}, & x_1 \neq x_2 \end{cases}$$