

Assignment 1 Consider the linear recursive specification

$$\begin{aligned} S &= a.T + \mathbf{1} \\ T &= b.S + b.U \\ U &= a.S + a.T + a.U \end{aligned}$$

- Construct, by transformation of the underlying recursive specification, a total and deterministic automaton that is language equivalent to the automaton $\mathcal{M}(S)$. (15 pts.)
- For $IA_\tau(\mathcal{N})$ we have the following result: If $x \approx y \cdot x + z$ then $x \approx y^* \cdot z$. Use the result to calculate an iteration expression for $\mathcal{L}(S)$.
(Hint: make use of deterministic equivalent of the given specification constructed for item a.)
(15 pts.)

Assignment 2

- Give the additional operational rules for the sequential operator \cdot and the iteration operator $*$ of the process algebra IA_τ , in total five rules. (5 pts.)
- Give the definition of (strong) bisimulation for processes of IA_τ . (5 pts.)
- Prove the law $x^* \Leftrightarrow x \cdot x^* + \mathbf{1}$ for IA_τ by showing a suitable relation R on IA_τ to be a bisimulation relation. Give an explicit definition of R . (15 pts.)

Assignment 3

- Formulate the Pumping Lemma for push-down languages. (10 pts.)
- Prove, using the Pumping Lemma for push-down languages, that the language L given by
$$L = \{ w \in \{a^n b^n c^n \mid n \geq 0\} \}$$
is not push-down. (15 pts.)

Assignment 4 Let the domain $D \subseteq \{0, 1, +\}^*$ be given by

$$D = \{ w+v \mid w, v \in \{0, 1\}^*, |w| = |v| \}$$

Construct a Turing machine that computes binary addition on D , i.e. the function on D that has the binary representation of $w + v$ as its output for input strings w and v , interpreted as binary numbers of an equal number of digits. (20 pts.)