

### Assignment 1

- a) Formulate the Pumping Lemma for regular languages. (10 pts.)
- b) Prove that the language  $L = \{ w \in \{a, b\}^* \mid \#_a(w) = \#_b(w) \}$  is *not* regular. (10 pts.)

### Assignment 2

- a) Give the operational rules of the iteration algebra  $IA_\tau(\mathcal{N})$  for the sequential composition and the iteration construct. (Five in total.) (10 pts.)
- b) Prove that  $x^* \cdot (x + \mathbf{1}) \stackrel{\text{bisim}}{\simeq} x^*$ , for all processes  $x \in IA_\tau(\mathcal{N})$ , by providing a bisimulation relation and verifying that it is a bisimulation relation indeed, or disprove that  $x^* \cdot (x + \mathbf{1}) \stackrel{\text{bisim}}{\simeq} x^*$  by providing a counter example and arguing why the respective processes cannot be bisimilar. (10 pts.)

### Assignment 3

- a) Draw a push-down automaton  $M$  (with no more than 3 states) that accepts the language  $L = \{ a^n b a^n \mid n \geq 0 \}$ . (10 pts.)
- b) Prove that  $L = \mathcal{L}(M)$ , either by verifying two set inclusions or by providing path invariants. (10 pts.)

**Assignment 4** Consider the following processes of the process algebra  $CA_\tau(\{P_1, P_2, Q_1, Q_2\})$ , that includes parallel composition, the encapsulation operator and the abstraction operator:

$$\begin{aligned} P_1 &= \ell!1.a_1.P_1 \\ P_2 &= \ell!2.a_2.P_2 \\ Q_1 &= \ell?1.Q_2 \\ Q_2 &= \ell?2.Q_1 \end{aligned}$$

- a) Give a derivation for the first transition of  $\tau_\ell(\partial_\ell(P_1 \parallel P_2 \parallel Q_1))$ . (10 pts.)
- b) Draw the transition system of  $\tau_\ell(\partial_\ell(P_1 \parallel P_2 \parallel Q_1))$  (10 pts.)

**Assignment 5** Provide a Turing machine  $M$  with no more than six states that is able to compare binary numbers of equal length, i.e. give a Turing machine  $M$  that returns on input of the string  $w > v$ , with  $w, v \in \{0, 1\}^*$  such that  $|w| = |v| > 0$ , the answer *yes* if the binary number represented by  $w$  is larger than the binary number represented by  $v$ , and returns the answer *no* otherwise. (20 pts.)