Real-Time Architectures
2003/2004

Mapping on execution platform

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Static organization

- Layering
- Architecture of Application: single executive or modular
- Architecture of RTOS: monolithic or micro kernel

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Mapping between layers

- Usually there are different units of concurrency in each layer
- The mapping must be made explicit
Single executive ("cyclic executive")

- Application is written as a single repetition, fast enough
  - to execute time driven tasks
  - to execute event driven tasks
- Interrupt service routines just record event occurrence (event flags)

- Advantages
  - needs little or no OS support
- Disadvantages:
  - lack of structure (not extensible)
    - spaghetti code, no clear task structure
  - repetition may grow to be inefficient
    - dependencies across several iterations
    - scalability: extension may destroy everything
  - cannot distribute

Single executive: example

```c
while true do
    if ItIsTimeFor (Keyboard) then /* check the keyboard */ fi;
    if ItIsTimeFor (Mouse) then /* check the mouse */ fi;
    if EventOccurred (TempSensor) then /* deal with event */ fi;
    WaitFor (Minimum time to next check or Event interrupt);
end
```

- Notes:
  - in fact simulation of concurrent task execution
  - sometimes the only choice (dependent on what the RTOS supports)
    - no RTOS, just a simple library suffices
  - the idea may recur even in a multi-tasking context
    - e.g. this may be the way that a sensor and a handler communicate

Tasks

- Introduced in the analysis phase
- Associated with handling an event

- Reasons for additional decomposition: design
  - (partly) sharing between tasks
    - e.g. (actions on) a shared data structure
  - enable tweaking individual subtasks in a processing pipeline
    - e.g. scalable parts
  - assigning relative importance to subtasks
    - less important processing steps are easily left out in case of overload
  - deal with temporal input bursts
    - e.g. separate event registration and actual processing
    - frequency differences
Task-related OS concepts

- **Process** ("program in execution")
  - defines a data space
  - concurrency transparency on single processor
  - has at least one associated thread
  - unit of distribution
  - unit of fault containment

- **Thread**
  - unit of concurrency
  - unit of scheduling
  - though with one thread per process, a process is often said to be the unit of scheduling
  - several threads 'live' within shared address space
  - i.e., in a process

Reasons to introduce processes

- Allow for distribution
- Introduce safety boundaries
- Follow solution structure
- (Increase concurrency level)
  - since more processes means more threads
  - hence, actually the purpose is to introduce more threads

Reasons to introduce threads

- Increase concurrency level
  - performance
  - hide latency
  - exploit platform concurrency
  - discriminate importance levels in activities
  - e.g. interrupt routines
- Deal with the natural concurrency
  - natural organization, structure
  - e.g. thread per event
  - thread per resource
  - thread per (active) external interaction sequence
Mapping of tasks

- Each task mapped onto a process
  - and, by result, onto a thread
  - then tasks can share resources only through the OS
  - or they approach them using message passing
    - the resource is maintained by a separate process

- Each task mapped onto a thread
  - newly created or repeatedly restarted

Multiplexing

- Several threads per task
  - then the actions of a task are distributed across several threads
  - a logical consequence of having active objects
    - active objects: data + operations + independent, private behavior .... “object with private thread”
    - simplest form: message passing process

- Several tasks per thread
  - efficiency concern in the mapping, e.g. when the tasks exclude each other

Consequences for scheduling

- Assume FPS (e.g. RMS) at task level
  - direct mapping task on thread/process
    - OS must support same policy
  - one thread executes several tasks completely
    - OS must support dynamic priority assignment
    - priority setting derived from event
  - one task is executed by several threads
    - the threads act as resources to be acquired by the task
    - dynamic priority assignment requirement
      - the priority of a thread is determined by the task it is working for
Consequences for scheduling

- Assume DPS (e.g. EDF) at task level
  - direct mapping task on thread/process
    - OS must support same policy
  - one thread executes several tasks completely
    - OS must support same policy
    - the deadline will do the trick
  - one task is executed by several threads
    - need to determine deadlines for subtasks
    - no clear strategy

Validation

- Need to have values for task execution times
  - dependent on platform, OS, compiler
  - direct measurement
  - profiling
- Strict requirements on OS
  - discussed later
  - need timing information of relevant system calls and OS operations
- Question:
  - what would portability mean in this context?
  - is a platform-independent real-time design possible?
    - what would be needed?
    - what would be the parameterization?
    - what are typical problems?