Visualization of Object-Oriented (Java) programs

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Abstract: We provide an explicit, consistent, execution model for OO programs, specifically Java, together with a tool that visualizes the model. This equips the student with a model to think and communicate about OO programs. Especially for an e-learning situation this is significant. Firstly, such a model supports the interaction with teachers and provides a sound basis for the understanding of other e-learning material concerning programming. Secondly, the tool supports acquiring proficiency with the model in an e-learning setting by giving precise information to hone his/her conceptualization of program execution. The model is advanced in that it accommodates multi-threading. The tool is flexible in that it directly uses the Java Virtual Machine for the visualization: extensions and adaptations can thus make use of the information the actual execution provides. A case in point is modeling the execution of code involving user interaction through the Graphic User Interface library. We consider several options here.

1 Introduction

Teaching programming has evolved substantially over the past thirty years, along at least three axes. Firstly, the languages have progressed from Pascal-like simple imperative languages that provide structuring through procedural decomposition to languages that are based on object-orientation (OO), which provide powerful additional structuring through object decomposition and inheritance between classes. Secondly, the use of library code has increased greatly; this, for example, makes programming Graphical User Interfaces (GUI’s) feasible. Thirdly, the role of programming in the curriculum has changed from being a typical nuts-and-bolts Computer Science subject to a more generally appreciated and applied skill that has a place in the curriculum of many departments.

This calls for new ways of offering tutoring in programming. Furthermore, to take advantage of what OO offers, the demanding concepts of this paradigm need to be well-understood.

From our experience in teaching programming (moving from Pascal, via C and C++ to Java) at varying departments at the Eindhoven University of Technology (The Netherlands) as well as in industry, we learned that an explicit conceptual execution model is indispensable, both for teacher - student communication as well as for communication between students. Indicative for the need for an execution model is that the paradigm is called object-orientation rather than class-orientation: thinking about the program is in terms of the objects that occur during execution rather than in terms of the classes that occur in the static code.

Especially for an e-learning situation students should have an explicit, consistent model to think and communicate about OO programs. Firstly, such a model supports the interaction with teachers and provides a sound basis for the understanding of other e-learning material concerning programming. Secondly, the tool supports acquiring proficiency with the model in an e-learning setting by giving precise information to hone his/her conceptualization of program execution.

In this paper we do not further argue the quite self-evident need for and use of such a model and tooling in e-learning, but concentrate on the model and visualization tool. We have implemented the ideas for Java, in a tool we named CoffeeDregs (in line with the Java/coffee association). Model and tool reflect our experience over the years; the current version is quite usable, but further evolution is envisaged. CoffeeDregs has a clear aim: it is a teaching tool that supports building a conceptual semantic model. It is not a debugger, nor does it visualize data or control for
diagnostic purposes as do the tools for professional programmers.

A key feature of the model is that it is detailed enough to explain the semantics of the language at statement level and abstract enough to enable thinking about the program at the level of objects. The distinction as well as the connection between dynamic execution in objects and static class code is made explicit. The emphasis is on stepwise modeling the dynamics of OO execution, i.e., the objects. Class code is merely present beside the visualization – model and tool indicate when and how the class code is involved in supplying information for the execution.

Another important feature is that execution of a program on the Java Virtual Machine (JVM) drives the visualization. This provides a close and direct connection to the actual execution of the language, enforcing the semantics to be realistic as well as facilitating extensions to the model and tool.

In this paper we give an overview of the model and tool; more detail can be found in (Luijten, 2003) and (Vandalon, 2010).

There are several approaches related to ours: Jeliot (Levy et al., 2002), Greenfoot (Kölling, 2010), BlueJ (Kölling et al., 2003), and JIVE (Gestwicki and Jayaraman, 2004; Lessa et al., 2010; Lessa and Jayaraman, 2012). Jeliot is more low-level than CoffeeDregs, for example showing in detail how expressions are evaluated. This makes Jeliot very well suited to the initial stages of programming education while CoffeeDregs also seeks to support more advanced concepts. Greenfoot and BlueJ enable to, visually, interact with instances of classes, but do not aim for modeling the detailed execution of a program – Greenfoot is especially well-suited for an informal explorative introduction. JIVE is most closely comparable to our approach. Also JIVE is much more ambitious in scope and application than CoffeeDregs: it is a dynamic analysis tool that aims to provide debugging facilities as well as being usable as a teaching tool. The lightweight approach in CoffeeDregs that restricts itself to “visualizing what happens inside the computer”, extended to complex concepts like user interaction through the Graphical User Interface. JIVE offers different, integrated views, for example using sequence diagrams to capture interactions over a longer time period. A technical difference is that JIVE is integrated with the IDE Eclipse, whereas CoffeeDregs is more IDE-independent. Also the use of Visual Operational Semantics (Jayaraman and Baltus, 1996)) and the Contour Model ((Johnston, 1971)) is subtly different. For a more detailed comparison and full treatment of the more subtle differences, see (Luijten, 2003).

We introduce our approach in section 2. In the next two sections we argue how to visualize, and what to show or hide. In section 5 we introduce the tool as applied to single threaded programs. In section 6 we briefly discuss its implementation. In section 7 we explain the extension to multi-threading. We consider options to treat user interaction through the GUI in section 8. This also provides an opportunity to explain our didactic approach in some more depth and shows the power of directly using information from the JVM for visualization.

2 Didactic approach

We use the bottom-up approach in teaching OO programming rather than the objects-first approach. One rationale is that we want to introduce the activity of a computer as performing small steps, for which it can be easily understood that they can be mechanized. Another reason is that, as our experience shows, it is advantageous to introduce objects not as given black boxes, but as composite entities explained in terms of already understood basic concepts. What we show initially is very simple: one class, one method, with basic statements. Then we incrementally add method decomposition and the OO concepts: object structures (references between objects, interaction through method calls), class structures (inheritance), library use, multi-threading, GUI programming, etc. The model shows progressively more of these concepts.

In our teaching we use lecture notes (we are aiming for a book version) with small chapters, introducing one concept at a time, providing its semantics in terms of the visualized, execution model. This semantics presents an abstract view of what happens inside the computer, e.g., variables and objects are visualized, stack frames are not. Thus we neither rely on machine notions for modeling nor do we make use of metaphor.

With each concept we provide a small example, the execution of which is visualized by the tool. The students then program more examples themselves, visualizing the execution with CoffeeDregs. Explaining the concepts in terms of the model together with a dynamic visualization proves effective in establishing a clear, concrete domain to think and communicate about programs. Note that with new concepts, new features are added to the model or already present ones are used in new ways, but that there is no change in the level of abstraction – only one model is built up incrementally.

CoffeeDregs can be used stand-alone, but also as a
plug-in to an IDE - we use the latter version for teaching, incorporated as a plug-in to NetBeans.

Choices about the model itself and how to visualize it are so much intertwined that from now on we will consider both under visualization.

3 How to visualize

The concepts that CoffeeDregs seeks to elucidate are the following.

1. Objects as container of data (variables) and center of manipulation (methods - with expressions, assignment); method and object recursion.
2. Scoping and life-time.
3. Object structures: references between objects and interaction through method calls.
4. Difference between class and object.
5. Class structures: inheritance (notably overriding of methods).
6. Use of libraries.
7. Multi-threading.
8. User interaction through GUI.

Essentially, the execution is visualized as a series of snapshots of the “inside of the computer” that shows the OO structure as well as the changes during execution. Guided by the above aims, the model has the following features (see Figure 1).

The guiding idea is to visualize the execution as a series of snapshots of the “inside of the computer” that shows the OO structure as well as the points of activity and the changes caused by this activity during execution. The lower limit level of abstraction is variable values, the abstraction level of the dynamics is stepwise statement execution. One snapshot at the time is visible, the effect of a step is indicated by the change in the snapshot, with additional clues like highlighting the locus of change.

Expression evaluation occurs as one step if only values or variables are involved. If a function is used in an expression, its execution involves the steps conform the visualization of method execution. Including detailed expression evaluation is one of the extensions of the tool that is considered.

1. We use a modified version of the Visual Operational Semantics (VOS) for OO by Jayarman and Baltus (Jayaraman and Baltus, 1996), which in turn is based on the Countour Model (CM) for scoping by Johnston (Johnston, 1971) (for more detail about motivation and use, see (Luijten, 2003)). We first and foremost model the dynamics of the execution and therefore only use contours, boxes, for objects and, inside these, for variable values and methods (also inside methods we use boxes to hold the values of local variables). As we explain the semantics at statement level, inside the contours the corresponding program code is present. Practicalities come to the fore here: to make this fit on the screen, rather than the full code of the body body, only a window of three lines of code (the active statement plus its immediate surroundings) may be displayed at each snapshot. A cursor indicates the place of control.

2. We use the contours to indicate scope, and also to indicate life-time: objects are visualized on creation - automated garbage collection is modeled in that objects that are out of scope are no longer displayed. Method calls are visualized as the appearance of the called method; methods are only displayed as long as they are active. Various subtleties apply here: for example, a shaded cursor at the place of a method call indicates that control (that moved to the executing method) will return there after the execution is finished. Note, that method recursion is automatically taken care of: each recursive call causes a new copy of the method being added to the object.

3. Object structures: references between objects are provided both as addresses (following the actual language mechanism) and through visible pointers that make the structure easier to comprehend. The visualization of method call and corresponding active method shows the interaction between objects, emphasizing that calls come from the outside, but that the activity itself is taking place inside the object.

4. The difference between class and object is made clear by separately displaying the class code, as code - not using contours. The objects are displayed in the main window on the screen, code is displayed to the right in a separate window. We indicate by a colored bar in the static code where this code is used to update the dynamic part: for example, when a new object is created or when a method is called.

NB Static variables and methods, notably main are incorporated in the model using a special class-object. As the treatment is similar to the regular objects, we do not further elaborate here.

5. Class structures: inheritance is modeled in the dynamic part by labeling objects with their dynamic type. A colored bar in the static part shows the
method that is selected - the code is copied to the dynamic part. In case of overriding, the appropriate method in the appropriate class is selected - directly: stepping through the inheritance hierarchy is not visualized. This is another example of the many subtle choices that have to be made: Visualizing the stepping, for example by moving the colored bar through the class code, could be a viable alternative, but might also create confusion about what a step in an execution means, and compromise the idea that class code is static. Therefore, for now we have chosen to limit visualizing activity inside objects. Note that variables in the object may come from different levels of the inheritance hierarchy. For simplicity, we do not consider variable shadowing, although the VCS can support this.

6. Use of library code is in principle covered by the model - to make the visualization feasible more needs to be done: see section 4.

7. Multi-threading is in principle straightforward: adding a thread identifier to the control-cursor: see section 7.

8. User-interaction through GUI is a challenging extension: see section 8.

**Remark**

CoffeeDregs only visualizes what happens inside the computer during execution. Therefore, input and output are not part of the visualization: the visualization is driven by the actual execution of the program, so input and output are no different from the case where the program is executed without visualization. In teaching this proves to be quite helpful in that the students do not get confused about what belongs to the visualization and what does not.

## 4 What to visualize

To make a visualization practicable requires careful consideration of what to display and what not to display. When executing a Java program, many auxiliary objects are created, resulting in a very large and incomprehensible visualization. Based on experience with previous versions of the tool, we investigated in detail how to decide which objects and methods to show and how to responsibly hide the others.

### 4.1 Ordering to visualization importance

We define a set of rules in terms of dynamic object properties and previously assigned object order. The rules are applied in order for each object. The first-matching rule decides the order of the object.

1. A method is currently executing in the object, assign order *method-executing*.
2. A method is active in the object, assign order *methods-active*.
3. the object is user-selected, assign order *user-selected*.
4. The object has no active methods and the object is referenced by at least 1 object of higher order than *methods-active*, assign order *object-referenced-from-above*.
5. The object is referenced by x other objects, assign order *object-often-referenced*.
6. No rule matches, assign order *bottom*.

The result is that objects with active methods get a high order, while inactive objects get lower orders. If an inactive object is directly referenced by an active object or if it is heavily referenced, it gets a slightly higher order.

We distinguish three basic levels of visibility.

1. Expanded. These contours show the full information, including values of instance variables and methods that are currently active in the object.
2. Collapsed. These contours show only the type and the reference of the object and no inner structure or text.
3. Hidden. These objects are not visualized on the screen.

The importance level is mapped to the visualization level as follows.

- Objects of order *methods-active* and higher are expanded.
- Objects of order between *object-often-referenced* and *object-referenced-from-above* are collapsed.
- Objects of order *bottom* are hidden.

### 4.2 Tying objects back together

All objects have an importance assigned to them and as a result some of them now have become hidden. Now other higher-importance objects might have become dangling because they were only referenced from currently hidden objects. We replace reference connections through objects of lower importance with transitive reference connections. They differ in appearance from the normal object references.

If there is a reference path from an expanded object A (see Figure 2), via one collapsed object S and
Figure 1: CoffeeDregs in action
then via one or more hidden objects to another expanded object \(B\), then there a transitive reference from \(S\) to \(B\) is added. If there is a direct reference from a collapsed object \(S\) to an expanded object \(B\) (i.e. there are zero hidden objects in between them), then there is a normal reference from \(S\) to \(B\).

### 4.3 Discussion

One could think of additional sublevels of expanded objects, in which for instance private variables are hidden if the object inherits from a non user-defined class. Beware that much of the functionality lies in these private variables and that, for example, a storage class can not show how the user-defined objects are stored within it if its private variables are hidden! The hiding of private variables should therefore probably occur for “leaves” in the object reference graph.

Great care is therefore required not to confuse accessibility in OO with reachability (and thus visibility) in the model. The aim of the model is to visualize the state of the program and not to visualize the structures in the program code. As a result, private variables may be visible, even though they are inaccessible from the program code. They nevertheless contribute to the execution model that is visualized and are therefore required to be visible.

If it is really necessary to hide variables, a better way probably is to merge them into a single `super-variable` \textit{private} which holds all references and thus the reference connections in the graph of the object model.

### 5 CoffeeDregs

Visualization is covered for the first six concepts described in section 3, in the manner described in section 4.

Programs are developed in an IDE, like the NetBeans we currently use. CoffeeDregs can be used stand-alone, taking the class files as well as the compiled program as input. Alternatively, CoffeeDregs can be used as a plug-in for NetBeans: CoffeeDregs is started from the NetBeans menu by clicking the CoffeeDregs icon in the menu.

Stepping through the program is push-button.

It is possible to step in reverse direction, e.g., to better understand what led to certain behavior. Currently there is no support for going back and changing the course of the execution by, e.g., providing different input or changing values of variables. From a didactic point of view it is debatable whether it would be desirable to have this option: on the one hand pursuing alternatives may help to understand the behavior of a program, on the other hand it might introduce confusion between interaction with a program and interaction with its visualization. For a technical motivation see section 6.

### 6 Implementing the visualization

The visualization and the program under study run in different virtual machines (JVMs). Using the Java Debug Interface, events are sent from the second JVM to the first when the program makes an execution step and these are used to visualize the state and to pause and resume execution of the second JVM. Reverse stepping is implemented by caching previous states.

The state is visualized using the Visual Object and Execution Model from the NetBeans Visual Library.

We lightly explore the software architecture. For more information on this topic, see (Vandalon, 2009)

CoffeeDregs is made up of two main components thereby partly following the Model-View-Controller pattern: debugmodel and debugview. The debugmodel handles the communication with the Java debugger and keeps a consistent state model. The debugview makes a drawing of the state model and acts as a controller to debugmodel.

When a program is loaded into CoffeeDregs, a VM is started for the program. CoffeeDregs then connects to the VM to subscribe to method entry and exit events and to execution step events. These events are reused as events within CoffeeDregs to notify the visualization of updates in the state.
When the visualization receives a notification, it updates its set of visualized objects. New objects in the programmer’s program are added to the visualization. Old objects that do not exist anymore in the programmer’s program are also removed from the visualization.

After the set of visualized objects has been updated, the objects in the set are updated to reflect changed values. If a method is active in an object, it is also added or its state is updated.

Following the update of the state is a reevaluation of the objects’ properties and applying the rules as described in section 4.

7 Multi-threading in CoffeeDregs

Multi-threaded programs are covered as follows: for each thread a separate control cursor is visualized, labeled with a number. This works well for the programs that occur in educational context, that typically have only a limited number of threads. This is specially instructive for Swing (GUI) programs where the Swing thread is visibly different from the main thread. Whether thread objects are visualized is optional.

There is no control as yet over which thread to advance.

8 User interaction through GUI

Modern languages like Java provide extensive support for user interaction, by means of libraries and a sophisticated framework to process user events.

We have experienced that explicitly treating (an abstract version of) the event loop is crucial in teaching this framework and prevents or takes away many misconceptions. For example, to explain the handling of mouse events, we show code as presented in figures 4 and 5.

Since this abstract code is not part of a Java implementation, CoffeeDregs can not visualize it. CoffeeDregs currently gives three options to visualize the corresponding Swing implementation.

One option is to hide all objects and activities from library code. This is the default. This gives the simplest view, but completely hides the event loop.

The second option is to visualize a few selected objects involved in the event loop in a collapsed view. This way, one sees when control is at the objects implementing the event loop, but the execution itself will not be visible.

The third option is to expand the objects that implement the event loop. CoffeeDregs will then give a detailed view of the execution of the event loop. This may be instructive, although its level of detail and hairiness may be detrimental to the understanding of many students.

The subject of current research is to see whether these possibilities suffice to support the explanation of the event model or that is feasible to extend CoffeeDregs with the option to visualize the abstract event loop.

9 Conclusions and future work

We have presented a viable visualization tool, CoffeeDregs, for OO (Java) programs. In its present form, the tool supports multi-threaded Java programs, with standard input and output and GUI-programs to a limited extent.

At this stage of the development of the tool, it is important to assess its educational effectiveness with quantitative experiments. Some preliminary steps in qualitative assessment have been reported on in (Luijten, 2003)

A new direction of development would be to connect the tool to a verification tool. Programs then carry annotations for verification, and the verification tool verifies whether or not the program satisfies these annotation. Having the visualization available, this would provide opportunities to better understand where and why annotations might fail. Promising candidate verification tool are Dafny (Leino, 2010) and Cocktail (Franssen, 2000).

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Control moves from `run` to `dispatchEvent` to the matching handler of the `eventListener`.

```java
class SwingThread {
    EventQueue&lt;AWT event&gt; eventQueue; //filled by the system
    AWTEvent e;
    void run {
        while (true) {
            if (!eventQueue.isEmpty()) {
                e = eventQueue.getNextEvent(); //take topmost event from eventQueue
                e.getSource().dispatchEvent(e); //on component that is source, call dispatchEvent
            }
        }
    }
}
```

Figure 4: Abstract event loop presented separately from CoffeeDregs

For each source that is a (sub)class (of) `JComponent`, lists of references to listeners and a `dispatchEvent` method are provided as follows.

```java
class JComponent {
    List&lt;MouseListener&gt; mouseListeners; //references to listeners registered on component,
    .../one list for each listener type
    void dispatchEvent(Event e) {
        if (isMouseListenerEvent(e)) { //for e MouseListenerEvent,
            for (MouseListener lis : mouseListeners) { //for each MouseListener registered on component
                if (isMouseClickedEvent(e)) { //depending on type of e,
                    lis.mouseClicked(e);
                } else if (isMousePressedEvent(e)){
                    lis.mousePressedEvent(e);
                } else if
                    ...//remaining MouseListenerEvent cases
            }
        } else if
            ...//remaining Event types
    }
}
```

Figure 5: Dispatch method in event source, presented separately from CoffeeDregs
REFERENCES


