

Models of Computation: Automata and Processes

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Motivation

- Automata theory provides simple models of computation for understanding the principles of computing and analysis of *computability*.
- Process theory has its origins in automata theory but focuses more on studying the notion of *interaction* and parallel behaviour.
- Goal: the *integration* of automata and process theory.
- The attempt at integration will reveal differences and similarities. We can use *analogies* between the theories to make the integration explicit.
- Add process theory to the undergraduate curriculum.

Automata and Process Theory: Similarity and Differences

Automata and Equivalences

Automata accept a language (a set of sequences of symbols) as correct or wanted behaviour. An automaton can for example model a coffee-vending machine:

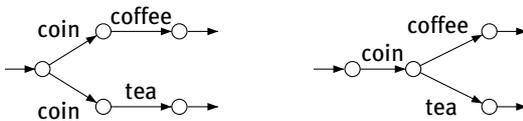


Figure 1: Two language equivalent automata

The above automata accept the same language, they are *language equivalent*:

- a coin followed by coffee
- a coin followed by tea

Process theory differentiates between them using *bisimulation equivalence*:

For a person using the machine it would make a difference whether inserting a coin predetermines the result or the choice is still available after inserting the coin.

Regular Expressions and Process Terms

- Regular expressions describe languages:
 $\text{coin} \cdot \text{coffee} + \text{coin} \cdot \text{tea}, \quad \text{coin} \cdot (\text{coffee} + \text{tea})$
- While regular expressions can describe all regular languages, their process term counterparts cannot describe all regular processes (shown in [1]).
- Process terms have calculation rules (axioms). E.g.:

$$(A3) \quad x + x = x$$

$$(A4) \quad (x + y)z = xz + yz$$

- The axiom $x(y + z) = xy + xz$ holds in automata theory but it does not hold in process theory!
- In process theory there are *additional operators*, such as \parallel , $|$, and $\llbracket _ \rrbracket$, for describing parallel behaviour which are not present in automata theory.

Grammars and Recursive Specifications

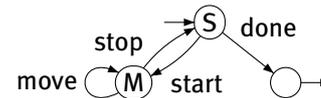


Figure 2: A regular process

- Grammars can also describe languages. The right-linear grammars from automata theory are equivalent to the recursive specifications of process theory.
- We can give both for the automaton in Figure 2:

$$\begin{array}{l|l} S \rightarrow \text{start } M \mid \text{done} & S = \text{start} \cdot M + \text{done} \\ M \rightarrow \text{move } M \mid \text{stop } S & M = \text{move} \cdot M + \text{stop} \cdot S \end{array}$$

Obtained Research Results

- In automata theory a context-free language can be accepted by an automaton using a stack (a pushdown automaton). In process theory, a context-free process can be transformed into a process communicating with the *Stack process*, making the interaction explicit. [2]
- Similar for basic parallel processes: a basic parallel process can be transformed into a process communicating with the *Bag process*. [3]
- Relative expressiveness between several classes has been investigated.

Extending the Chomsky Hierarchy

The Chomsky hierarchy discerns classes of languages (regular, context-free, etc.). The additional operators present in process theory create *new classes* of languages, such as the basic parallel class. The new classes create an extended, more fine-grained version of this hierarchy.

What does this new hierarchy look like? What can be expressed by each of these new classes? Do they have some finite axiomatisation?

Research Team

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References

- [1] J.C.M. Baeten, F. Corradini, and C.A. Grabmayer. A characterization of regular expressions under bisimulation. *Journal of the ACM*, 54(2), 2007.
- [2] J.C.M. Baeten, P.J.L. Cuijpers, and P.J.A. van Tilburg. A Context-Free Process as a Pushdown Automaton. *Proceedings of CONCUR'08, LNCS 5201*, pp. 98–113, 2008.
- [3] J.C.M. Baeten, P.J.L. Cuijpers, and P.J.A. van Tilburg. A Basic Parallel Process as a Parallel Pushdown Automaton. *Proceedings of EXPRESS'08, ENTCS*, 2008.