## Solutions for

## Tussentoets Wiskunde 2 (2DD52) Donderdag 27 November 2014, 11:00-12:30

Problem 1. We present a combined solution for both parts (a) and (b).

- We denote the four plant locations by their initial letters $G, L, M, N$, and we denote the three Soda types Blue, Green, Red by their initial letters $B, G, R$.
- For $i \in\{G, L, M, N\}$ and $j \in\{B, G, R\}$, we introduce a variable $x_{i j}$ that indicates the amount of soda of type $j$ (measured in $m^{3}$ ) to be produced at plant location $i$.
(1) The overall expected profit for soda of type $j$ produced at plant $i$ then equals the amount $x_{i j}$ multiplied by the expected profit per $m^{3}$ in $€$ for soda type $j$. Therefore the objective function is

$$
\begin{aligned}
\max Z= & 550\left(x_{G B}+x_{L B}+x_{M B}+x_{N B}\right) \\
+ & 400\left(x_{G G}+x_{L G}+x_{M G}+x_{N G}\right) \\
+ & 350\left(x_{G R}+x_{L R}+x_{M R}+x_{N R}\right)
\end{aligned}
$$

(2) At every plant $i$, the processed amount of water must not exceed the available quantity:

$$
\begin{aligned}
& x_{G B}+x_{G G}+x_{G R} \leq 400 \\
& x_{L B}+x_{L G}+x_{L R} \leq 900 \\
& x_{M B}+x_{M G}+x_{M R} \leq 550 \\
& x_{N B}+x_{N G}+x_{N R} \leq 350
\end{aligned}
$$

(3) At every plant $i$, the amount of worked hours must not exceed the available amount:

$$
\begin{aligned}
& 2 x_{G B}+4 x_{G G}+3 x_{G R} \leq 1800 \\
& 2 x_{L B}+4 x_{L G}+3 x_{L R} \leq 3200 \\
& 2 x_{M B}+4 x_{M G}+3 x_{M R} \leq 2300 \\
& 2 x_{N B}+4 x_{N G}+3 x_{N R} \leq 1200
\end{aligned}
$$

(4) For every soda type $j$, the total produced amount must not exceed the production ceiling:

$$
\begin{array}{r}
x_{G B}+x_{L B}+x_{M B}+x_{N B} \leq 700 \\
x_{G G}+x_{L G}+x_{M G}+x_{N G} \leq 800 \\
x_{G R}+x_{L R}+x_{M R}+x_{N R} \leq 300
\end{array}
$$

(5) Finally, all quantities $x_{i j}$ must be non-negative:

$$
x_{i j} \geq 0 \quad \text { for all } i \in\{G, L, M, N\} \text { and } j \in\{B, G, R\}
$$

## Problem 2.

(a) For making all variables non-negative, substitute $x_{2}^{\prime}=-x_{2}$ and $x_{3}^{\prime}=-x_{3}$.

For getting a maximization problem, multiply the objective function by -1 .
For getting non-negative right hand sides, multiply the inequalities by -1 .
For getting equality constraints, introduce slack and surplus variables $s_{1}, s_{2}, s_{3}$.
This yields the following linear program in standard form:

$$
\begin{array}{rllll}
\max W= & 3 x_{1} & +x_{2}^{\prime} & +x_{3}^{\prime} & \\
& & & \\
\text { s.t. } & -x_{1} & +x_{2}^{\prime} & +x_{3}^{\prime} & -s_{1} \\
& x_{1} & +x_{2}^{\prime} & & -s_{2} \\
& 2 x_{1} & -x_{2}^{\prime} & -x_{3}^{\prime} & \\
& & +s_{3} & =3 \\
& & & \\
\text { with } & x_{1}, x_{2}^{\prime}, x_{3}^{\prime}, s_{1}, s_{2}, s_{3} \geq 0 & &
\end{array}
$$

(b) For the first two constraints, the 2-phase method introduces artificial variables $a_{1}$ and
$a_{2}$. The objective in the first phase is to maximize $-a_{1}-a_{2}$.

|  | $W$ | $x_{1}$ | $x_{2}^{\prime}$ | $x_{3}^{\prime}$ | $s_{1}$ | $s_{2}$ | $s_{3}$ | $a_{1}$ | $a_{2}$ | $b$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $W$ | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 0 |
| $a_{1}$ | 0 | -1 | 1 | 1 | -1 | 0 | 0 | 1 | 0 | 1 |
| $a_{2}$ | 0 | 1 | 1 | 0 | 0 | -1 | 0 | 0 | 1 | 3 |
| $s_{3}$ | 0 | 2 | -1 | -1 | 0 | 0 | 1 | 0 | 0 | 2 |

The first row is not in the right shape, as some basic variables have non-zero cost coefficients. Hence we repair it in the following way:

|  | $W$ | $x_{1}$ | $x_{2}^{\prime}$ | $x_{3}^{\prime}$ | $s_{1}$ | $s_{2}$ | $s_{3}$ | $a_{1}$ | $a_{2}$ | $b$ |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $W$ | 1 | 0 | -2 | -1 | 1 | 1 | 0 | 0 | 0 | -4 |
| $a_{1}$ | 0 | -1 | 1 | 1 | -1 | 0 | 0 | 1 | 0 | 1 |
| $a_{2}$ | 0 | 1 | 1 | 0 | 0 | -1 | 0 | 0 | 1 | 3 |
| $s_{3}$ | 0 | 2 | -1 | -1 | 0 | 0 | 1 | 0 | 0 | 2 |

Variable $a_{1}$ leaves the basis, and $x_{2}^{\prime}$ enters:

|  | $W$ | $x_{1}$ | $x_{2}^{\prime}$ | $x_{3}^{\prime}$ | $s_{1}$ | $s_{2}$ | $s_{3}$ | $a_{1}$ | $a_{2}$ | $b$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $W$ | 1 | -2 | 0 | 1 | -1 | 1 | 0 | 2 | 0 | -2 |
| $x_{2}^{\prime}$ | 0 | -1 | 1 | 1 | -1 | 0 | 0 | 1 | 0 | 1 |
| $a_{2}$ | 0 | 2 | 0 | -1 | 1 | -1 | 0 | -1 | 1 | 2 |
| $s_{3}$ | 0 | 1 | 0 | 0 | -1 | 0 | 1 | 1 | 0 | 3 |

Variable $a_{2}$ leaves the basis, and $s_{1}$ enters:

|  | $W$ | $x_{1}$ | $x_{2}^{\prime}$ | $x_{3}^{\prime}$ | $s_{1}$ | $s_{2}$ | $s_{3}$ | $a_{1}$ | $a_{2}$ | $b$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $W$ | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 0 |
| $x_{2}^{\prime}$ | 0 | 1 | 1 | 0 | 0 | -1 | 0 | 0 | 1 | 3 |
| $s_{1}$ | 0 | 2 | 0 | -1 | 1 | -1 | 0 | -1 | 1 | 2 |
| $s_{3}$ | 0 | 3 | 0 | -1 | 0 | -1 | 1 | 0 | 1 | 5 |

The first phase terminates with $a_{1}=a_{2}=0$. We remove the columns for the artificial variables and start the second phase (with the original objective function):

|  | $W$ | $x_{1}$ | $x_{2}^{\prime}$ | $x_{3}^{\prime}$ | $s_{1}$ | $s_{2}$ | $s_{3}$ | $b$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $W$ | 1 | -3 | -1 | -1 | 0 | 0 | 0 | 0 |
| $x_{2}^{\prime}$ | 0 | 1 | 1 | 0 | 0 | -1 | 0 | 3 |
| $s_{1}$ | 0 | 2 | 0 | -1 | 1 | -1 | 0 | 2 |
| $s_{3}$ | 0 | 3 | 0 | -1 | 0 | -1 | 1 | 5 |

The first row is not in the right shape, as some basic variables have non-zero cost coefficients. Hence we repair it in the following way:

|  | $W$ | $x_{1}$ | $x_{2}^{\prime}$ | $x_{3}^{\prime}$ | $s_{1}$ | $s_{2}$ | $s_{3}$ | $b$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $W$ | 1 | -2 | 0 | -1 | 0 | -1 | 0 | 3 |
| $x_{2}^{\prime}$ | 0 | 1 | 1 | 0 | 0 | -1 | 0 | 3 |
| $s_{1}$ | 0 | 2 | 0 | -1 | 1 | -1 | 0 | 2 |
| $s_{3}$ | 0 | 3 | 0 | -1 | 0 | -1 | 1 | 5 |

Now we see that the reduced cost coefficient of $s_{2}$ is negative, while none of the other entries in the corresponding column is positive. We conclude that the LP is unbounded.
(c) The dual of (LP.1) looks as follows:

$$
\begin{array}{rlrll}
\max V=-y_{1} & -3 y_{2}-2 y_{3} \\
& & \\
\text { s.t. } & y_{1} & -y_{2} & -2 y_{3} & \leq-3 \\
& y_{1} & +y_{2} & -y_{3} & \geq 1 \\
& y_{1} & & -y_{3} & \geq 1
\end{array}
$$

with

$$
y_{1}, y_{2} \leq 0 ; y_{3} \geq 0
$$

(d) Since the primal (LP.1) is unbounded, its dual (LP.2) must be infeasible.

