

Reinder J. Brill, r.j.brill@tue.nl  
TU/e Informatica, System Architecture and Networking

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SAN

## Real-Time Architectures 2003/2004

Summary

Reinder J. Brill

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Reinder J. Brill, r.j.brill@tue.nl  
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## Summary

- Response times for FPPS
  - Without jitter
  - With release jitter
  - Completion (or end) jitter
  - Example
- Resource reservation

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## Response times for FPPS

Worst-case response times	Best-case response time
$WR_j = \sup_{\phi,i} R_{j,i}$	$BR_j = \inf_{\phi,i} R_{j,i}$
<i>Critical</i> instant	<i>Optimal</i> instant
$x = C_j + \sum_{i < j} \left\lceil \frac{x}{T_i} \right\rceil C_i$	$x = C_j + \sum_{i < j} \left( \left\lfloor \frac{x}{T_i} \right\rfloor - 1 \right) C_i$
<i>smallest</i> positive solution highest lower bound	<i>largest</i> positive solution lowest upper bound
solved by means of an iterative procedure	
can also be determined by means of a time line	

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### Response times with release jitter AJ

Worst-case response times	Best-case response time
$WR_j = \sup_{\phi, i} R_{j,i}$	$BR_j = \inf_{\phi, i} R_{j,i}$
<i>Critical instant*</i>	<i>Optimal instant*</i>
$x = C_j + \sum_{i < j} \left\lceil \frac{x + AJ_i}{T_i} \right\rceil C_i$	$x = C_j + \sum_{i < j} \left( \left\lfloor \frac{x - AJ_i}{T_i} \right\rfloor - 1 \right) C_i$
<i>smallest</i> positive solution highest lower bound	<i>largest</i> positive solution lowest upper bound
solved by means of an iterative procedure	
can also be determined by means of a time line	

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### Completion (or end) jitter EJ

- $EJ_j$  of a task  $\tau_j$  *excluding* release jitter:  

$$EJ_j = \sup_{\phi, i, k} (R_{j,i} - R_{j,k}) \leq WR_j - BR_j$$
- $EJ_j$  of a task  $\tau_j$  *including* release jitter:  

$$EJ_j \leq AJ_j + WR_j^* - BR_j^*$$
  - Where the \* denotes: including release jitter

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### Completion and release jitter

- Example:

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## Resource Reservation

- Definition:
  - Resource reservation is a technique (or *mechanism*) at the level of an RTOS, providing a *virtual platform*.
- Example:  $(T, B)$ ,  $T = 20$  ms,  $B = 5$  ms.
- Multiple resources:
  - (co-) processor(s), memory, bus, network, ...
- Motivation:
  - Priorities used to/for:
    - Guarantee deadlines (FPS and DPS);
    - Assign relative importance (FPS only);
    - Co-operation protocols + concurrency control.

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