



# Column Generation and its applications

Murat Firat, dept. IE&IS, TU/e  
BPI Cluster meeting  
March 14, 2018

**TU/e** Technische Universiteit  
Eindhoven  
University of Technology

**Where innovation starts**

# Outline

Some real-life decision problems

- Standard formulations

Basics of Column Generation

- Master formulations

Case: Shift scheduling of airport workers

- Problem description

- Master formulation

- Reduced cost and pricing problem

Column Generation overview

Towards an integer solution

Using data science in the CG approach

# Outline

Some real-life decision problems

Standard formulations

Basics of Column Generation

Master formulations

Case: Shift scheduling of airport workers

Problem description

Master formulation

Reduced cost and pricing problem

Column Generation overview

Towards an integer solution

Using data science in the CG approach

# Different industry sectors, different problems

- ▶ Telecommunication: Stable workforce assignments.

# Different industry sectors, different problems

- ▶ Telecommunication: Stable workforce assignments.
- ▶ Telecommunication: Designing FTTH network.

# Different industry sectors, different problems

- ▶ Telecommunication: Stable workforce assignments.
- ▶ Telecommunication: Designing FTTH network.
- ▶ Airport ground operations: Shift scheduling of workers.

# Different industry sectors, different problems

- ▶ Telecommunication: Stable workforce assignments.
- ▶ Telecommunication: Designing FTTH network.
- ▶ Airport ground operations: Shift scheduling of workers.
- ▶ Logistics: Planning routes of vehicles.

# Different industry sectors, different problems

- ▶ Telecommunication: Stable workforce assignments.
- ▶ Telecommunication: Designing FTTH network.
- ▶ Airport ground operations: Shift scheduling of workers.
- ▶ Logistics: Planning routes of vehicles.
- ▶ Machine learning: Constructing max-accuracy decision trees.



# Instance sizes

- ▶ Stable workforce assignments: 100 workers, 25 jobs, 15 skills.

# Instance sizes

- ▶ Stable workforce assignments: 100 workers, 25 jobs, 15 skills.
- ▶ Designing FTTH network: Local areas of 5-10K citizens,

# Instance sizes

- ▶ Stable workforce assignments: 100 workers, 25 jobs, 15 skills.
- ▶ Designing FTTH network: Local areas of 5-10K citizens,
- ▶ Shift scheduling of airport workers: 200 workers, 4 weeks, 5 skills.

## Instance sizes

- ▶ Stable workforce assignments: 100 workers, 25 jobs, 15 skills.
- ▶ Designing FTTH network: Local areas of 5-10K citizens,
- ▶ Shift scheduling of airport workers: 200 workers, 4 weeks, 5 skills.
- ▶ Planning routes of vehicles: 100-1K customers, 50 vehicles.

# Instance sizes

- ▶ Stable workforce assignments: 100 workers, 25 jobs, 15 skills.
- ▶ Designing FTTH network: Local areas of 5-10K citizens,
- ▶ Shift scheduling of airport workers: 200 workers, 4 weeks, 5 skills.
- ▶ Planning routes of vehicles: 100-1K customers, 50 vehicles.
- ▶ Decision trees: 5K data instances, 50 attributes.

# Common properties of the problems

- ▶ Strict feasibility requirements.

# Common properties of the problems

- ▶ Strict feasibility requirements.
- ▶ Worst cases: NP-Hard.

# Common properties of the problems

- ▶ Strict feasibility requirements.
- ▶ Worst cases: NP-Hard.
- ▶ Exponentially many feasible solutions



# Common properties of the problems

- ▶ Strict feasibility requirements.
- ▶ Worst cases: NP-Hard.
- ▶ Exponentially many feasible solutions
- ▶ Many local optimal points.

# Standard formulations

One big formulation, low-quality bound.

- ▶ Stable workforce assignments: Worker-job decisions.

# Standard formulations

One big formulation, low-quality bound.

- ▶ Stable workforce assignments: Worker-job decisions.
- ▶ Shift scheduling of airport workers: Worker-shift decisions.

# Standard formulations

One big formulation, low-quality bound.

- ▶ Stable workforce assignments: Worker-job decisions.
- ▶ Shift scheduling of airport workers: Worker-shift decisions.
- ▶ Planning routes of vehicles: Vehicle-customer decisions.

# Standard formulations

One big formulation, low-quality bound.

- ▶ Stable workforce assignments: Worker-job decisions.
- ▶ Shift scheduling of airport workers: Worker-shift decisions.
- ▶ Planning routes of vehicles: Vehicle-customer decisions.
- ▶ Decision trees: Data row-tree node decisions.

# Outline

Some real-life decision problems

Standard formulations

Basics of Column Generation

Master formulations

Case: Shift scheduling of airport workers

Problem description

Master formulation

Reduced cost and pricing problem

Column Generation overview

Towards an integer solution

Using data science in the CG approach

# Master formulations

- ▶ Stable workforce assignments: **Team**-job decisions.

# Master formulations

- ▶ Stable workforce assignments: **Team**-job decisions.
- ▶ Shift scheduling of airport workers: Worker-**schedule** decisions.



# Master formulations

- ▶ Stable workforce assignments: **Team**-job decisions.
- ▶ Shift scheduling of airport workers: Worker-**schedule** decisions.
- ▶ Planning routes of vehicles: Vehicle-**route** decisions.

# Master formulations

- ▶ Stable workforce assignments: **Team**-job decisions.
- ▶ Shift scheduling of airport workers: **Worker-schedule** decisions.
- ▶ Planning routes of vehicles: **Vehicle-route** decisions.
- ▶ Decision trees: **Data segment**-tree node decisions.

# Master formulation

- ▶ Linear programming model:

# Master formulation

- ▶ Linear programming model:
  - More complicated objects as decision variables.

# Master formulation

- ▶ Linear programming model:
  - More complicated objects as decision variables.
- ▶ Object set

# Master formulation

- ▶ Linear programming model:
  - More complicated objects as decision variables.
- ▶ Object set
  - has exponentially many items.

# Master formulation

- ▶ Linear programming model:
  - More complicated objects as decision variables.
- ▶ Object set
  - has exponentially many items.
  - initially is empty or has few items.

# Master formulation

- ▶ Linear programming model:
  - More complicated objects as decision variables.
- ▶ Object set
  - has exponentially many items.
  - initially is empty or has few items.
- ▶ Iteratively find promising columns and add them to object set.



# Master formulation

- ▶ Linear programming model:
  - More complicated objects as decision variables.
- ▶ Object set
  - has exponentially many items.
  - initially is empty or has few items.
- ▶ Iteratively find promising columns and add them to object set.
- ▶ When no promising column exists: certificate for optimality.

# Main loop of Column Generation

- ▶ Reduced cost of a column is

# Main loop of Column Generation

- ▶ Reduced cost of a column is
  - violation amount of the corresponding dual constraint.

# Main loop of Column Generation

- ▶ Reduced cost of a column is
  - violation amount of the corresponding dual constraint.
  - estimated cost change per unit increase in its value.

# Main loop of Column Generation

- ▶ Reduced cost of a column is
  - violation amount of the corresponding dual constraint.
  - estimated cost change per unit increase in its value.
- ▶ Adding promising columns

# Main loop of Column Generation

- ▶ Reduced cost of a column is
  - violation amount of the corresponding dual constraint.
  - estimated cost change per unit increase in its value.
- ▶ Adding promising columns
  - completes dual feasibility.

# Main loop of Column Generation

- ▶ Reduced cost of a column is
  - violation amount of the corresponding dual constraint.
  - estimated cost change per unit increase in its value.
- ▶ Adding promising columns
  - completes dual feasibility.
  - improves primal objective.

# Main loop of Column Generation

- ▶ Reduced cost of a column is
  - violation amount of the corresponding dual constraint.
  - estimated cost change per unit increase in its value.
- ▶ Adding promising columns
  - completes dual feasibility.
  - improves primal objective.
- ▶ No dual feasibility improvement  $\approx$  primal optimality



# Outline

Some real-life decision problems

Standard formulations

Basics of Column Generation

Master formulations

Case: Shift scheduling of airport workers

Problem description

Master formulation

Reduced cost and pricing problem

Column Generation overview

Towards an integer solution

Using data science in the CG approach

# Problem definition: Shift scheduling

We are given:

- ▶ multi-skilled workers with availability info.

# Problem definition: Shift scheduling

We are given:

- ▶ multi-skilled workers with availability info.
- ▶ service demand within a planning horizon

# Problem definition: Shift scheduling

We are given:

- ▶ multi-skilled workers with availability info.
- ▶ service demand within a planning horizon
- ▶ labor regulations about shifts:

# Problem definition: Shift scheduling

We are given:

- ▶ multi-skilled workers with availability info.
- ▶ service demand within a planning horizon
- ▶ labor regulations about shifts:
  - minimum resting time between shifts,

# Problem definition: Shift scheduling

We are given:

- ▶ multi-skilled workers with availability info.
- ▶ service demand within a planning horizon
- ▶ labor regulations about shifts:
  - minimum resting time between shifts,
  - maximum working time due to contracts,

# Problem definition: Shift scheduling

We are given:

- ▶ multi-skilled workers with availability info.
- ▶ service demand within a planning horizon
- ▶ labor regulations about shifts:
  - minimum resting time between shifts,
  - maximum working time due to contracts,
  - night shifts: longer resting times.

# Basic notation

- ▶ Let  $\mathcal{S}$  be set of all schedules st. a schedule  $s \in \mathcal{S}$



# Basic notation

- ▶ Let  $\mathcal{S}$  be set of all schedules st. a schedule  $s \in \mathcal{S}$ 
  - covers time interval  $i \in N$  if  $s^i = 1$ , not otherwise.

# Basic notation

- ▶ Let  $\mathcal{S}$  be set of all schedules st. a schedule  $s \in \mathcal{S}$ 
  - covers time interval  $i \in N$  if  $s^i = 1$ , not otherwise.
  - should comply all regulations

# Basic notation

- ▶ Let  $\mathcal{S}$  be set of all schedules st. a schedule  $s \in \mathcal{S}$ 
  - covers time interval  $i \in N$  if  $s^i = 1$ , not otherwise.
  - should comply all regulations
  - should have length bet. 3-9 hrs.

# Basic notation

- ▶ Let  $\mathcal{S}$  be set of all schedules st. a schedule  $s \in \mathcal{S}$ 
  - covers time interval  $i \in N$  if  $s^i = 1$ , not otherwise.
  - should comply all regulations
  - should have length bet. 3-9 hrs.
- ▶ For worker  $w \in W$

# Basic notation

- ▶ Let  $\mathcal{S}$  be set of all schedules st. a schedule  $s \in \mathcal{S}$ 
  - covers time interval  $i \in N$  if  $s^i = 1$ , not otherwise.
  - should comply all regulations
  - should have length bet. 3-9 hrs.
- ▶ For worker  $w \in W$ 
  - $Sk_w^d$  indicates if worker  $w$  is skilled in skill type  $d \in D$ .

# Basic notation

- ▶ Let  $\mathcal{S}$  be set of all schedules st. a schedule  $s \in \mathcal{S}$ 
  - covers time interval  $i \in N$  if  $s^i = 1$ , not otherwise.
  - should comply all regulations
  - should have length bet. 3-9 hrs.
- ▶ For worker  $w \in W$ 
  - $Sk_w^d$  indicates if worker  $w$  is skilled in skill type  $d \in D$ .
- ▶ Service demand  $R_{i,d}$  at time  $i$  in skill type  $d$ ,

# Master IP formulation

$$\begin{aligned} \text{Min} \quad & \sum_{w \in W} \sum_{s \in \mathcal{S}} c_{sw} x_{sw} \\ \text{subject to} \quad & \sum_{w \in W} S k_w^d s^i x_{sw} \geq R_{i,d}, \quad d \in D, i \in N \\ & \sum_{s \in \mathcal{S}} x_{sw} = 1, \quad w \in W \\ & x_{sw} \in \{0, 1\}, \quad s \in \mathcal{S}, w \in W \end{aligned}$$

# LP Relaxation of Master formulation

$$\begin{aligned} & \text{Min} && \sum_{w \in W} \sum_{s \in S'} c_{sw} x_{sw} \\ & \text{subject to} && \\ & && \sum_{w \in W} S k_w^d s^i x_{sw} \geq R_{i,d}, \quad d \in D, i \in N \\ & && \sum_{s \in S'} x_{sw} = 1, \quad w \in W \\ & && 0 \leq x_{sw} \leq 1, \quad s \in S', w \in W \end{aligned}$$

Note: Restricted set  $S' \subset S$



# LP Relaxation of Master formulation

$$\begin{array}{ll} \text{Min} & \sum_{w \in W} \sum_{s \in S'} c_{sw} x_{sw} \\ \text{subject to} & \\ & \sum_{w \in W} S k_w^d s^i x_{sw} \geq R_{i,d}, \quad d \in D, i \in N \\ & \sum_{s \in S'} x_{sw} = 1, \quad w \in W \\ & 0 \leq x_{sw} \leq 1, \quad s \in S', w \in W \end{array} \quad \begin{array}{l} \text{Duals} \\ \pi_{i,d} \\ \theta_w \end{array}$$

Note: Restricted set  $S' \subset S$

## Reduced cost of column $x_{sw}$

Dual constraint of column  $x_{sw}$ :

$$\sum_{d \in D} \sum_{i \in N} S k_w^d s^i \pi_{i,d}^* + \theta_w^* \leq c_{sw} \quad (1)$$

Reduced cost of column  $x_{sw}$ :

$$\bar{c}_{sw} = c_{sw} - \sum_{d \in D} \sum_{i \in N} S k_w^d s^i \pi_{i,d}^* - \theta_w^* \quad (2)$$

**Case**  $\bar{c}_{sw} < 0$ : (1) Estimated objective decrease (**why?**), (2) dual feasibility violation.

## Pricing problem: Objective

Pricing problem: Find the most promising column (schedule) with the objective

$$\min_{s \in \mathcal{S}, w \in \mathcal{W}} \left\{ c_{sw} - \sum_{d \in D} \sum_{i \in N} S k_w^d s^i \pi_{i,d}^* - \theta_w^* \right\} \quad (3)$$

$$= \max_{s \in \mathcal{S}, w \in \mathcal{W}} \left\{ \sum_{i \in N} s^i \left( \sum_{d \in D} S k_w^d \pi_{i,d}^* - c_{iw} \right) \right\} - \theta_w^* \quad (4)$$

# Pricing problem: Modeling

Define a graph,



- ▶ reds: shifts, blues: resting times, blacks: unavailability.

# Pricing problem: Modeling

Define a graph,



- ▶ reds: shifts, blues: resting times, blacks: unavailability.
- ▶ labor regulations: arc structure and side constraints

# Pricing problem: Modeling

Define a graph,



- ▶ reds: shifts, blues: resting times, blacks: unavailability.
- ▶ labor regulations: arc structure and side constraints
- ▶ Solve  $\max_a w_a x_a$  subject to

# Pricing problem: Modeling

Define a graph,



- ▶ reds: shifts, blues: resting times, blacks: unavailability.
- ▶ labor regulations: arc structure and side constraints
- ▶ Solve  $\max_a w_a x_a$  subject to
  - conservations, resting after shifts, passing unavailability arcs.

# Pricing problem: Modeling

Define a graph,



- ▶ reds: shifts, blues: resting times, blacks: unavailability.
- ▶ labor regulations: arc structure and side constraints
- ▶ Solve  $\max_a w_a x_a$  subject to
  - conservations, resting after shifts, passing unavailability arcs.
- ▶ **Pricing**: find the constrained  $0 - |N|$  "Longest Path"!



# Pricing problem: Modeling

A schedule  $s$  on the graph looks like:



# Column Generation overview

- ▶ Initialization:

# Column Generation overview

- ▶ Initialization:
  - Formulate Master ILP model, obtain the Restricted Master Problem (RMP).
  - Express the reduced costs, formulate pricing.

# Column Generation overview

- ▶ Initialization:
  - Formulate Master ILP model, obtain the Restricted Master Problem (RMP).
  - Express the reduced costs, formulate pricing.
- ▶ **Warm up:** Find several initial columns for a warm start of RMP.

# Column Generation overview

- ▶ Initialization:
  - Formulate Master ILP model, obtain the Restricted Master Problem (RMP).
  - Express the reduced costs, formulate pricing.
- ▶ **Warm up:** Find several initial columns for a warm start of RMP.
- ▶ **Step 1:** Solve the Restricted Master Problem, pass duals to pricing.

# Column Generation overview

- ▶ Initialization:
  - Formulate Master ILP model, obtain the Restricted Master Problem (RMP).
  - Express the reduced costs, formulate pricing.
- ▶ **Warm up:** Find several initial columns for a warm start of RMP.
- ▶ **Step 1:** Solve the Restricted Master Problem, pass duals to pricing.
- ▶ **Step 2:** Solve pricing:

# Column Generation overview

- ▶ Initialization:
  - Formulate Master ILP model, obtain the Restricted Master Problem (RMP).
  - Express the reduced costs, formulate pricing.
- ▶ **Warm up:** Find several initial columns for a warm start of RMP.
- ▶ **Step 1:** Solve the Restricted Master Problem, pass duals to pricing.
- ▶ **Step 2:** Solve pricing:
  - If  $\exists i : \bar{c}_i < 0$ : update  $\mathcal{S}'$ , go to **Step 1**.
  - If  $\forall i : \bar{c}_i \geq 0$ : RMP is solved to optimality, go to **Step 3**.

# Column Generation overview

- ▶ Initialization:
  - Formulate Master ILP model, obtain the Restricted Master Problem (RMP).
  - Express the reduced costs, formulate pricing.
- ▶ **Warm up:** Find several initial columns for a warm start of RMP.
- ▶ **Step 1:** Solve the Restricted Master Problem, pass duals to pricing.
- ▶ **Step 2:** Solve pricing:
  - If  $\exists i : \bar{c}_i < 0$ : update  $\mathcal{S}'$ , go to **Step 1**.
  - If  $\forall i : \bar{c}_i \geq 0$ : RMP is solved to optimality, go to **Step 3**.
- ▶ **Step 3:** Output the RMP solution.



# Outline

Some real-life decision problems

Standard formulations

Basics of Column Generation

Master formulations

Case: Shift scheduling of airport workers

Problem description

Master formulation

Reduced cost and pricing problem

Column Generation overview

Towards an integer solution

Using data science in the CG approach

# Towards integer solutions

Having fractional RMP optimal solution, we have two choices:

- ▶ Use rounding heuristics:

# Towards integer solutions

Having fractional RMP optimal solution, we have two choices:

- ▶ Use rounding heuristics:
  - Use meta heuristics to find a feasible solution quickly (hopefully)

# Towards integer solutions

Having fractional RMP optimal solution, we have two choices:

- ▶ Use rounding heuristics:
  - Use meta heuristics to find a feasible solution quickly (hopefully)
  - **Make decisions** how to (smartly) round the fractional solution

# Towards integer solutions

Having fractional RMP optimal solution, we have two choices:

- ▶ Use rounding heuristics:
  - Use meta heuristics to find a feasible solution quickly (hopefully)
  - **Make decisions** how to (smartly) round the fractional solution
- ▶ Start a smart enumeration, e.g. Branch-and-Price, either

# Towards integer solutions

Having fractional RMP optimal solution, we have two choices:

- ▶ Use rounding heuristics:
  - Use meta heuristics to find a feasible solution quickly (hopefully)
  - **Make decisions** how to (smartly) round the fractional solution
- ▶ Start a smart enumeration, e.g. Branch-and-Price, either
  - to obtain "optimal" integer sol'n.
  - to output "best-found" integer sol'n with quality measure in a time limit.

# Outline

Some real-life decision problems

Standard formulations

Basics of Column Generation

Master formulations

Case: Shift scheduling of airport workers

Problem description

Master formulation

Reduced cost and pricing problem

Column Generation overview

Towards an integer solution

Using data science in the CG approach

# Using data science: Advantage or waste?

Consider a data of instance history and their solutions.

- ▶ Find "similar" instances and round the current fractional solution towards corresponding solutions.



# Using data science: Advantage or waste?

Consider a data of instance history and their solutions.

- ▶ Find "similar" instances and round the current fractional solution towards corresponding solutions.
- ▶ Simply find "similar" previous solutions to the current fractional solution and round accordingly.

# Using data science: Advantage or waste?

Consider a data of instance history and their solutions.

- ▶ Find "similar" instances and round the current fractional solution towards corresponding solutions.
- ▶ Simply find "similar" previous solutions to the current fractional solution and round accordingly.
- ▶ Cluster instances and analyze the commonalities (patterns) in the cluster solutions.

THANKS